Universality of distances in random graphs

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YEP workshop, March 10-14, 2008

Joint work with:

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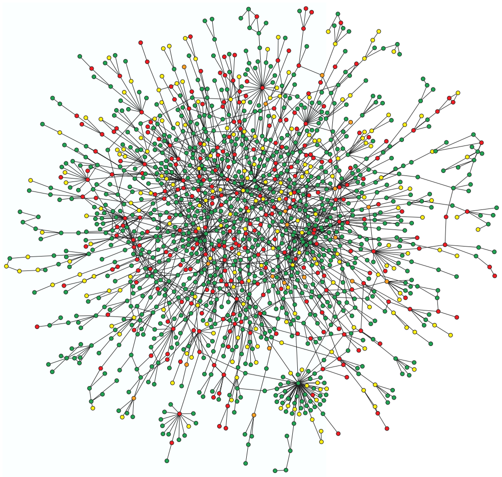
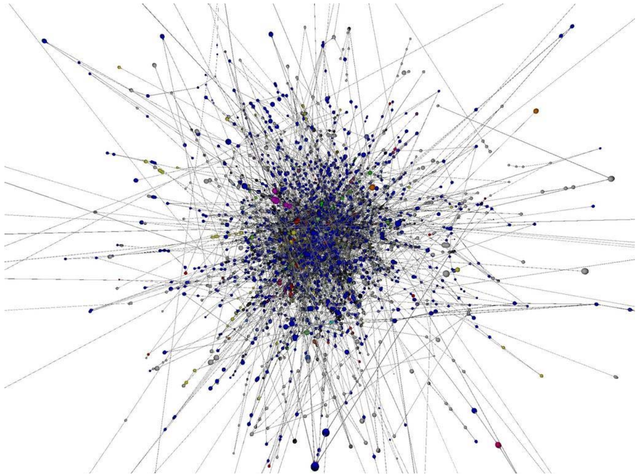
- Piet van Mieghem (TUD)
- Dmitri Znamenski (EURANDOM, now Philips Research)

Complex networks

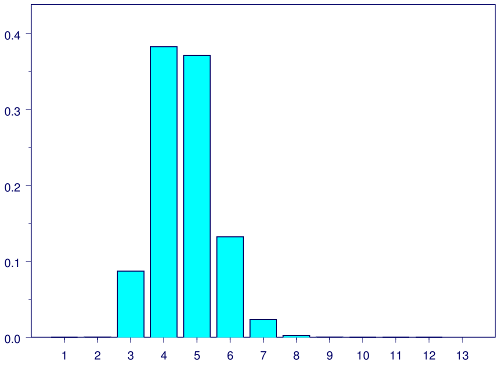


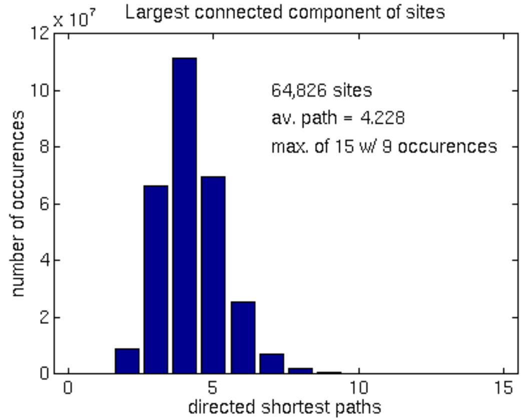
Yeast protein interaction network



Internet topology in 2001

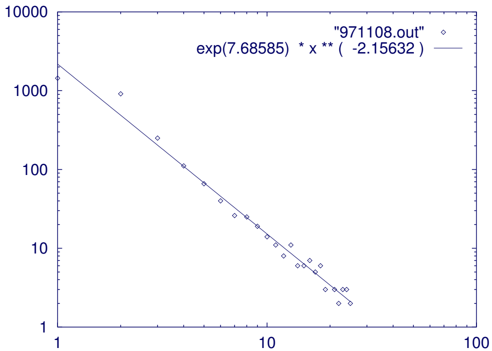
Small-world phenomenon

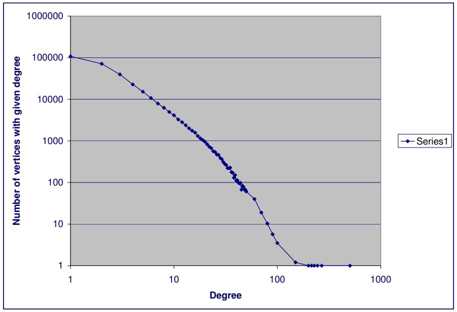




Distances in AS graph and WWW (Adamic 99)

Scale-free phenomenon





Loglog plot of degree sequences in AS graph in Internet in 1997 (FFF97) and in the collaboration graph among mathematicians (`http://www.oakland.edu/enp`)

Modeling complex networks

- Inhomogeneous Random Graphs:
Static random graph, independent edges with inhomogeneous edge occupation probabilities, yielding scale-free graphs.
(BJR07, CL02, CL03, BDM-L05, CL06, NR06, EHH06,...)

- Configuration Model:
Static random graph with prescribed degree sequence.
(MR95, MR98, RN04, HHV05, EHHZ06, HHZ07, JL07, FR07,...)

- Preferential Attachment Model:
Dynamic random graph, attachment proportional to degree plus constant.
(BA99, BRST01, BR03, BR04, M05, B07, HH07,...)

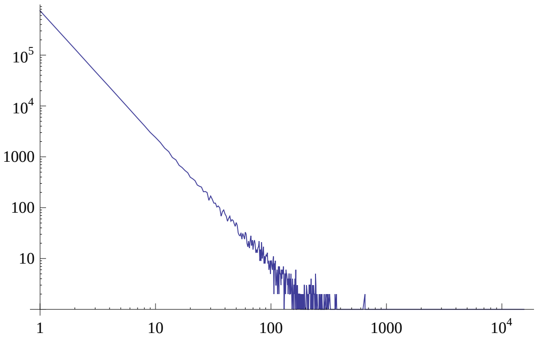
Configuration model

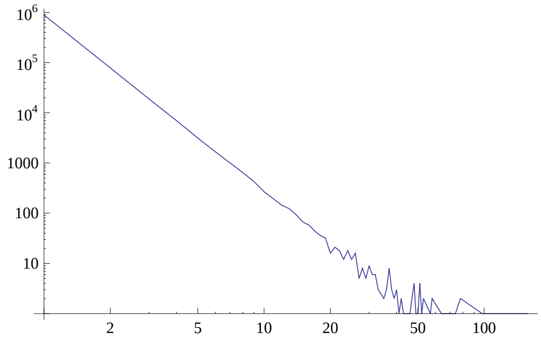
N is number of vertices. Consider i.i.d. sequence of degrees $D_1, D_2, \ldots, D_N$, with

$$\mathbb{P}(D_1 \geq k) = c_\tau k^{-\tau+1}(1 + o(1)),$$

where c_τ is normalizing constant and $\tau > 1$.

Power law degree sequence CM





Loglog plot of degree sequence CM with $n = 1.000.000$ and $\tau = 2.5$ and $\tau = 3.5$, respectively.

Configuration model: graph construction

How to construct graph with above degree sequence?

- Assign to vertex j degree D_j.

$$L_N = \sum_{i=1}^{N} D_i$$

is total degree. Assume L_N is even.
Incident to vertex i have D_i 'stubs' or half edges.

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• Connect stubs to create edges as follows:
Number stubs from 1 to L_N in any order.
First connect first stub at random with one of *other* $L_N - 1$ stubs.
Continue with second stub (when not connected to first) and so on, until all stubs are connected...

Distances in configuration model

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$$H_N \approx \log_\nu N, \qquad \text{with} \qquad \nu = \frac{\mathbb{E}[D(D-1)]}{\mathbb{E}[D]}.$$

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- When $\tau \in (1,2)$, (EHHZ06)

$$H_N \quad \text{uniformly bounded.}$$

Consequences and proof

Proof relies on coupling of neighborhood of vertices to branching process.

Extensions:

• Fluctuations around leading order are uniformly bounded, and we compute its 'limiting distribution'.

• Diameter of graph is maximal distance between any pair of connected vertices.

Diameter CM of order $\log N$ when $\mathbb{P}(D_i \geq 3) < 1$ (FR07),
while of order $\log \log N$ when $\tau \in (2, 3)$ and $\mathbb{P}(D_i \geq 3) = 1$ (HHZ07).

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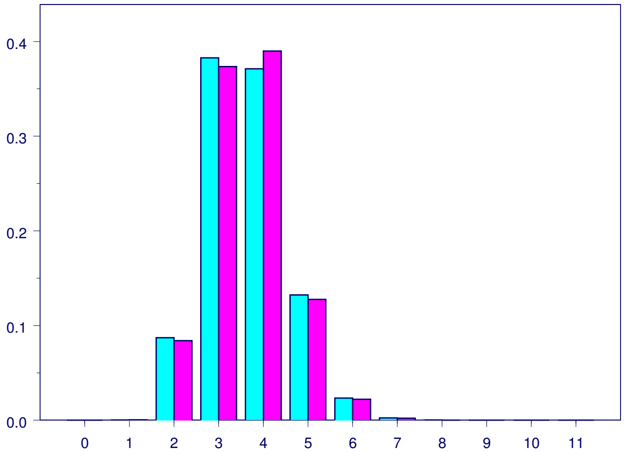
• Similar methodology applies to various inhomogeneous random graphs, with similar (but generally weaker) results.

Comparison Internet data



Number of AS traversed in hopcount data (blue) compared to the model (purple) with $\tau = 2.25, N = 10,940.$

Preferential attachment

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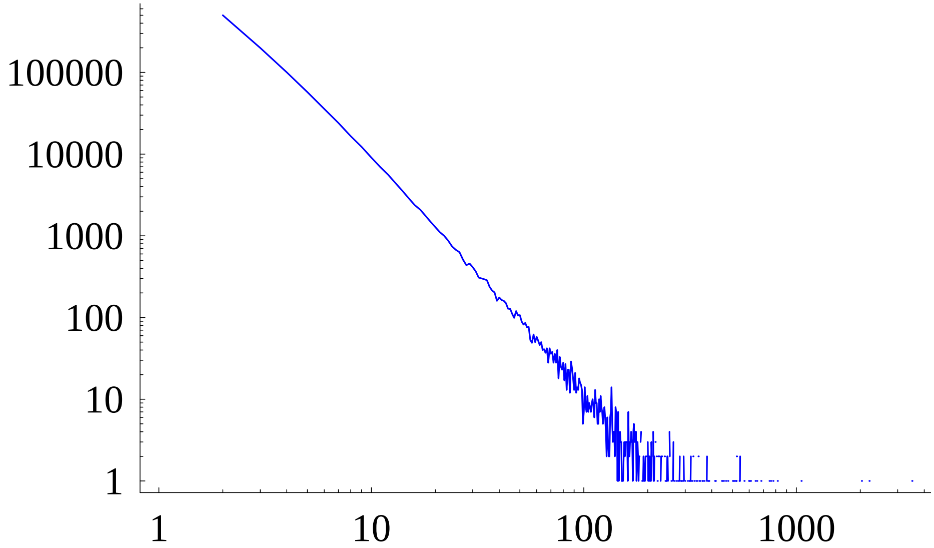
• Different edges can attach with different updating rules:
(a) intermediate updating degrees with self-loops (BA99, BR04, BRST01)
(b) intermediate updating degrees without self-loops;
(c) without intermediate updating degrees. i.e., independently.

(Graphs in cases (b-c) have advantage of being connected.)

Scale-free nature PA

Yields power-law degree sequence with exponent $\tau = 3 + \delta/m \in (2, \infty)$.



$$(m = 2, \delta = 0, \tau = 3 + \frac{\delta}{m} = 3)$$

Distances PA models

- Diameter PA bounded above by $C \log t$ for all $m \geq 1$ and $\delta > -m$ (b-c) (HH07);

- Diameter PA bounded above by $C \log \log t$ for $m \geq 2$ and $\delta \in (-m, 0)$ for which power-law exponent $\tau \in (2,3)$ (a-c) (HH07);

- Diameter PA bounded above by $(1+\varepsilon)\frac{\log t}{\log \log t}$ for $m \geq 2$ and $\delta = 0$ for which power-law exponent $\tau = 3$ (a) (BR04);

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- Diameter PA bounded below by $\varepsilon \log t$ for $m \geq 2$ and $\delta > 0$ for which power-law exponent $\tau \in (3, \infty)$ (a-c) (HH07).

Universality PA models

First evidence of strong form of universality:
random graphs with similar degree structure share similar behavior.

For random graphs, universality predicted by physics community.

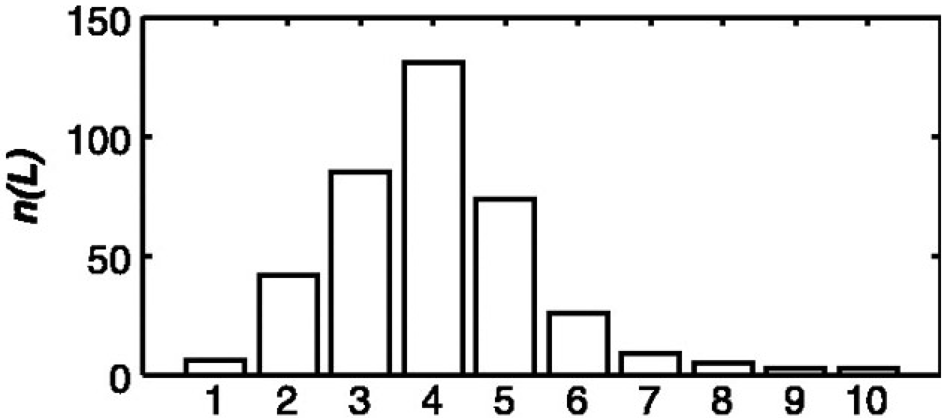
Universality is leading paradigm in statistical physics.
Only few examples where universality can be rigorously proved.

Key question: Can universality be proved for processes such as Ising model or contact process on random graphs?

More information on random graphs:

`www.win.tue.nl/∼rhofstad/NotesRGCN.pdf`

Small-world phenomenon



Distances in social network (Small-World Project Watts (2003))