

Phase transition for the vacant set left by a random walk on
random d -regular graph

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Eurandom, 16 December 2009

Model.

- ▶ $G_n = (V_n, \mathcal{E}_n)$ – sequence of finite connected graphs, $|V_n| \rightarrow \infty$.
- ▶ $(X_t^n)_{t \geq 0}$ – cont. time RW on G_n started from its invariant distribution
- ▶ $u > 0$ – parameter
- ▶ $\mathcal{V}^u = V_n \setminus \{X_t : t \in [0, u|V_n|]\}$ – vacant set.

Question: Percolative properties of the random set \mathcal{V}^u .

Is there $u_c = u_c(G) \in (0, \infty)$ such that

- ▶ for $u > u_c$, \mathcal{V}^u has only ‘small’ components (say $O(\log n)$).
- ▶ for $u < u_c$, \mathcal{V}^u has a ‘giant’ component (a component with volume $O(|V_n|)$).
- ▶ is the giant component unique
- ▶ how does \mathcal{V}^u look at/close to u_c .

The problem was studied only on d -**dimensional torus** $(\mathbb{Z}/n\mathbb{Z})^d$, $d \geq 3$:

- ▶ Benjamini-Sznitman (JEMS '08): If u is small enough, then \mathcal{V}^u has a giant component. This component is unique in some weak sense.
- ▶ improved slightly by D. Windisch (EJP '08)
- ▶ D. Windisch (ECP '08): Local picture on the torus converges to *Random interlacement* on \mathbb{Z}^d .
- ▶ Recent results on Random Interlacement can be used to prove:
 - ▶ For $u < u_{**} < u_*(\mathbb{Z}^d)$, there is a giant component
 - ▶ For $u > u_{***} > u_*(\mathbb{Z}^d)$, there are only small components.

Here $u_*(\mathbb{Z}^d)$ is the percolation threshold for Random Interlacement.

- ▶ It is believed that $u_c((\mathbb{Z}/n\mathbb{Z})^d) = u_*(\mathbb{Z}^d)$.

Our setting

Consider graphs that are simpler for Bernoulli percolation:

- ▶ random d -regular graph
(graph uniformly chosen from all d -regular graphs on n vertices)
- ▶ d -regular large-girth expanders

Both these classes of graphs are *good finite approximations of d -regular tree*

Bernoulli percolation on such graphs studied by Alon-Benjamini-Stacey '04, Nachmias-Peres '09, Pittel '09.

More precisely, we consider graphs satisfying:

(A0) $G_n = (V_n, \mathcal{E}_n)$ is d -regular, $|V_n| = n$.

(A1) *Local almost tree-like property:*

There exists $\alpha_1 \in (0, 1)$ such that for all n and $x \in V_n$

the ball $B(x, \alpha_1 \log n)$ contains at most one cycle

(A2) *Uniform spectral gap:*

There exists $\alpha_2 > 0$ such that for all n

$$\lambda_1(G_n) \geq \alpha_2$$

Remarks

- ▶ random d -regular graph satisfies (A0)–(A2) with high probability as $n \rightarrow \infty$.
- ▶ (A1) \implies typical point $x \in V_n$ has *tree-like* neighbourhood of radius $\frac{\alpha_1}{2} \log n$.
- ▶ (A2) is equivalent (via Cheeger's inequality) to expansion property.

$$\frac{|\partial A|}{|A|} \geq \alpha'_2, \quad \forall n, \forall A \subset V_n, |A| < |V_n|/2.$$

Theorem

Let G_n satisfy (A0)–(A2). Then there exists explicitly computable constant $u_c = u_c(d)$ such that

1. when $u > u_c$, there is $K = K(u, d, \alpha_1, \alpha_2)$ such that

$$\mathbb{P}[|\mathcal{C}_{\max}(\mathcal{V}^u)| \geq K \log n] \xrightarrow{n \rightarrow \infty} 0.$$

2. when $u < u_c$, there is $\rho = \rho(u, d, \alpha_1, \alpha_2)$ such that

$$\mathbb{P}[|\mathcal{C}_{\max}(\mathcal{V}^u)| \geq \rho n] \xrightarrow{n \rightarrow \infty} 1.$$

3. when $u < u_c$, for every $\varepsilon > 0$,

$$\mathbb{P}[|\mathcal{C}_{\text{sec}}(\mathcal{V}^u)| \geq \varepsilon n] \xrightarrow{n \rightarrow \infty} 0.$$

Here, $\mathcal{C}_{\max}(\mathcal{V}^u)$ ($\mathcal{C}_{\text{sec}}(\mathcal{V}^u)$) is the largest (second largest) component of \mathcal{V}^u .

Remark. 3. is equivalent with: $\exists f_n = o(n)$ such that

$$\mathbb{P}[|\mathcal{C}_{\text{sec}}(\mathcal{V}^u)| \geq f_n] \xrightarrow{n \rightarrow \infty} 0.$$

We are far from $f_n = O(\log n)$.

Relation to Random Interlacement model

Let $\mathbb{G} = (\mathbb{V}, \mathcal{E}_{\mathbb{G}})$ be infinite, transient, connected graph, which is local limit of G_n :
 $\exists r_n \rightarrow \infty$ such that for typical $x \in V_n$, the ball $B_{G_n}(x, r_n) \stackrel{\phi_n^x}{\simeq} B_{\mathbb{G}}(\emptyset, r_n)$.

Examples: Torus $\rightarrow \mathbb{Z}^d$, our $G_n \rightarrow d$ -regular tree \mathbb{T}_d .

Random Interlacement is a dependent percolation on \mathbb{G} which is local limit of \mathcal{V}^u (should be proved).

Local definition of RI: Let $A \subset \mathbb{V}$ finite.

- ▶ *equilibrium measure:* $e_A(x) = \mathbb{1}_A(x) P_x[\text{RW never returns to } A]$.
- ▶ at every point x start $\text{Poisson}(ue_A(x))$ independent random walks
- ▶ $\text{RI}|_A$ has the same law as the set of vertices in A not visited by these random walks.
- ▶ *notation:* \mathbb{C}_x cluster of RI containing x .

Conjecture. For $u_*(\mathbb{G}) = \sup\{u : \mathbb{P}[|\mathbb{C}_x| = \infty] > 0\}$, the threshold of RI on \mathbb{G} ,

$$u_c(G) = u_*(\mathbb{G}).$$

Random Interlacement on the tree has a particularly simple description in terms of branching process (Teixeira '09):

Lemma

Given $\emptyset \in RI$, the cluster \mathbb{C}_x has the law of branching process whose offspring distribution is binomial with parameters $d - 1$ (resp. d in the first generation) and p_u , where

$$p_u = \exp \left\{ - \frac{u(d-2)^2}{d(d-1)} \right\}.$$

Consequence. RI on \mathbb{T}_d exhibits a phase transition on $u_*(\mathbb{T}_d)$ given by $p_{u_*(\mathbb{T}_d)}(d-1) = 1$.

Theorem

If G_n satisfy (A0)–(A2), then

$$u_c(G) = u_*(\mathbb{T}_d).$$

Local convergence to Random Interlacement

RI on the tree \mathbb{T}_d is indeed a good local model for the vacant set on G_n :

Lemma

There is $\beta \in (0, \alpha_1/5)$, such that for all x with tree-like neighbourhood of radius $5\beta \log n$, for all $u > 0$, $\varepsilon > 0$, there exists a coupling \mathbb{P} of RW on G and RI's on \mathbb{T}^d such that

$$\mathbb{P}[\mathbb{C}_\emptyset^{u-\varepsilon} \supset \phi_n^x(\mathcal{C}_x(\mathcal{V}^u \cap B(x, \beta \log n))) \supset \mathbb{C}_\emptyset^{u+\varepsilon}] \xrightarrow{n \rightarrow \infty} 1.$$

Consequence. In every tree like ball of radius $\beta \log n$ we have a good control of $\mathcal{C}_x(\mathcal{V}^u)$ by a branching process.

THE LOCAL CONTROL IS NOT SUFFICIENT!

- ▶ In the super-critical phase, the giant component cannot be contained in a ball of radius $\beta \log n$.
- ▶ In the sub-critical phase, the largest cluster is $\sim K \log n$, but $K \rightarrow \infty$ as $u \downarrow u_c$. (In particular, since $\text{diam } G = \log n(1 + o(1))$, we have $|\mathcal{C}_{\max}(\mathcal{V}^u)| \geq \text{diam } G$.)

Segments-Bridges measure

Random walk is not suitable for our techniques (there is a dependence between its positions at various times).

We replace the RW trajectory by a different trajectory which possesses more independence and approximates RW well.

Segments-Bridges construction.

- ▶ Fix $L = n^\gamma$, $\ell = \log^2 n$, where $\gamma \in (0, 1)$. Set $m = un/(L + \ell)$.
- ▶ Let $(Y^i)_{i < m}$ be an i.i.d. family with values in $D([0, L], V)$ and marginal P_π^L (RW started from invariant distribution, killed after L steps) = **segments**
- ▶ given (Y^i) , let $(Z^i)_{i < m}$ (**bridges**) be independent family of $D([0, \ell], V)$ valued random variables,

$$\text{law of } Z^i = P_{Y_L^i, Y_0^{i+1}}^\ell = \text{RW bridge of length } \ell \text{ from } Y_L^i \text{ to } Y_0^{i+1}$$

Lemma

Law of concatenation $Y^0 Z^0 Y^1 \dots Y^m Z^m$ approximates well $(X_t : t \leq un)$.

Proof: Due to spectral gap assumption, the mixing time of the RW is $O(\log n)$.

Sub-critical regime

We want to show: For some $K \geq 0$,

$$P[\mathcal{C}_{\max}(\mathcal{V}^u) \geq K \log n] \xrightarrow{n \rightarrow \infty} 0.$$

Here $\mathcal{V}^u = V \setminus \cup_{i < m} (Y^i \cup Z^i)$.

We are looking for an upper bound. It is sufficient to control $\bar{\mathcal{V}}^u = V \setminus \cup_{i < m} Y^i$:

$$\mathbb{P}[|\mathcal{C}_x(\bar{\mathcal{V}}^u)| \geq K \log n] \leq n^{-1-\varepsilon}.$$

Construct $\mathcal{C}_x(\bar{\mathcal{V}}^u)$ sequentially via a breath-first-search algorithm.

PICTURE!

The algorithm discovers some information on segments Y^i :

- ▶ \mathcal{A}_k – information discovered during exploration of first $k - 1$ vertices.
- ▶ y_k , k -th explored vertex.

To control the algorithm one need to control the probability that the vertex y_k is vacant given \mathcal{A}_k .

$$\mathbb{P}[y_k \in \bar{\mathcal{V}}^u | \mathcal{A}_k] < \frac{1}{d-1}.$$

Control of the BFS algorithm

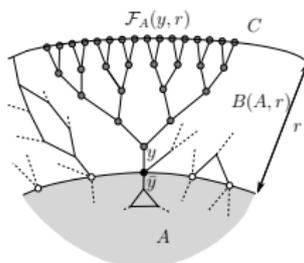
We have

$$\mathbb{P}[y_k \in \bar{\mathcal{V}}^u | \mathcal{A}_k] = \mathbb{P}[y_k \notin \cup_{i < m} Y^i | \mathcal{A}_k] \leq \mathbb{P}[y_k \notin \cup_{i \text{ non-marked}} Y^i | \mathcal{A}_k].$$

On \mathcal{A}_k , $(Y^i)_{i \text{ non-marked}}$ are i.i.d. with the distribution $P_{\pi}^L[\cdot | H_{A_k} > L]$, where $A_k = \{y_1, \dots, y_{k-1}\}$ is the set of already explored points.

We need to compute $P[H_{y_k} > L | H_{A_k} > L] = \frac{P[H_{y_k \cup A_k} > L]}{P[H_{A_k} > L]}$.

This can be done, but only in a specific situation ($r = 7 \ln \ln n$):



Why:

- ▶ By Aldous-Brown (1993) $P[H_A > L] \sim e^{-L/\mathbb{E}[H_A]}$
- ▶ $\mathbb{E}[H_A] \sim C_{\text{eff}}(A, B(A, r))^{-1}$
- ▶ We need to control $C_{\text{eff}}(A, B(A, r)) - C_{\text{eff}}(A \cup y, B(A, r))$.

Deterministic estimate on bad event

In this specific situation one gets for non-marked Y^i ,

$$\mathbb{P}[y_k \notin Y^i | \mathcal{A}_k] = 1 - Ln^{-1}(1 + o(1)) = 1 - n^{\gamma-1}(1 + o(1)).$$

and thus, when $u > u_c$,

$$\mathbb{P}[y_k \in \bar{\mathcal{V}}^u | \mathcal{A}_k] = \mathbb{P}[y_k \notin \cup_i \text{non-marked } Y^i | \mathcal{A}_k] = (1 - n^{\gamma-1})^m < 1/(d-1).$$

Recall, $m = un/(L + \ell) \sim un^{1-\gamma}$.

One further needs:

- ▶ the set of marked segments is much smaller m .
- ▶ (A1) \implies The good situation does not occur at most crK^2 times during the run of the algorithm (deterministically)

Strategy for Bernoulli percolation. $p > p_c$

1. Fix $p' \in (p_c, p)$.
2. Use local branching-process control to construct many components of size n^β , $\beta \in (0, 1)$:

$$\mathbb{P}_{p'}[|\{x : |\mathcal{C}_x| \geq n^\beta\}| \geq \rho n] \xrightarrow{n \rightarrow \infty} 1.$$

3. Sprinkle additional edges over G (increase from p' to p).
These edges will with high probability joint the components constructed in 2.
(local \rightarrow global control)
4. For 3, the expansion property is very useful.
(It implies $\forall A, B \subset V$, $|A|, |B| \geq cn$, there exist $\Theta(n)$ disjoint paths from A to B all having length smaller than C .)

It is not clear how to implement the sprinkling for the random walk trajectory!

First idea of sprinkling

1. In super-critical regime $u < u_c$. So, fix $u' \in (u, u_c)$.
2. Consider $\mathcal{V}^{u'}$ and use the local branching-process comparison to construct many large components:

$$\mathbb{P}[|\{x : |\mathcal{C}_x(\mathcal{V}^{u'})| \geq n^\beta\}| \geq \rho n] \xrightarrow{n \rightarrow \infty} 1.$$

3. Erase the last part of the trajectory, that is $\{X_t : t \in [un, u'n]\}$.
4. BUT! What is the law of this part given $\mathcal{V}^{u'}$??
5. We were not able to control it.

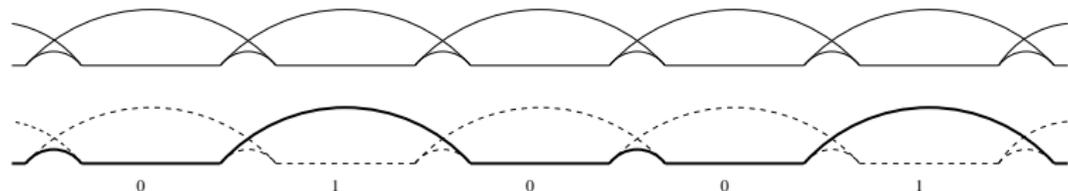
Second idea of sprinkling

1. Use Segment-Bridges measure as before for $u' \in (u, u_c)$. That is set $m' = u'n/(L + \ell)$
2. By local branching-process comparison, construct many large components as before.
3. Consider sequence $(R_i)_{i < m'}$ of i.i.d. Bernoulli random variables with parameter $(u' - u)/u'$, independent of Y^i 's and Z^i 's.
4. Erase segments Y^i with $R^i = 1$.
5. PROBLEM: We obtain non-nearest-neighbour path.

Solution. Put even more bridges into the picture, so that erasing some segments does not disconnect the trajectory:

Long range bridges Z^{ij} , $i < m'$, $j < \log n$.

$$\text{law of } Z^{ij} = P_{Y_L^i, Y_0^{i+j}}^\ell$$



Technical complications

The union of segments and all long range bridges is not n.n. path, but quite complicated object \implies it is hard to obtain the local control by branching process.

1. Establish control by the branching process for complement of segments only.
2. Use this control to prove:

$$\mathbb{P}\left[|\{x \in V : x \text{ is proper in } \bar{\mathcal{V}}^{u'}\}| \geq cn\right] \xrightarrow{n \rightarrow \infty} 1,$$

where x is proper if

- ▶ $|\mathcal{C}_x(\bar{\mathcal{V}}^{u'} \cap B(x, \beta \log n))| \geq c((d-1)p_{u'+\varepsilon})^{\beta \log n}$.
 - ▶ $|\mathcal{C}_x(\bar{\mathcal{V}}^{u'}) \cap \partial B(x, k)| \leq ((d-1)p_{u'-\varepsilon})^k$ for all $k \in [c' \log \log n, \beta \log n]$.
3. Show that if number of proper points is larger than cn , then flipping $n^{1-\gamma} \log^5 n$ vertices does not destroy a proportion of large components (deterministic statement):
After adding long range bridges:

$$|\{x : \mathcal{C}_x \geq c((d-1)p_{u'+\varepsilon})^{\beta \log n}\}| \geq cn/8.$$

The proof follows the same lines as the proof of existence of giant component.

Sprinkling shows that any two giant components at level u' join at level u .

Problem: the sprinkling can also create new giant component at level u , that was not there at level u'

It is actually not possible, since any giant component at level u must contain a positive proportion of vertices that were in n^β components at level u' .

By sprinkling again, we know that these are joined together.

1. Density of giant cluster for $u < u_c$.
2. Stronger uniqueness result. Is $|\mathcal{C}_{\text{sec}}| = O(\log n)$.
3. What about $u = u_c$. Is there the double jump as in the Bernoulli case:
 $|\mathcal{C}_{\text{max}}| \sim n^{2/3}$? Critical window?

The end

Thank you for your attention.