

Metastates and short-range spin glasses

Part II(b)

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Ground states and Domain walls

Definition of ground state

Domain walls

Domain walls for ground states

Domain wall measure

Ground state domain wall geometry

Definitions

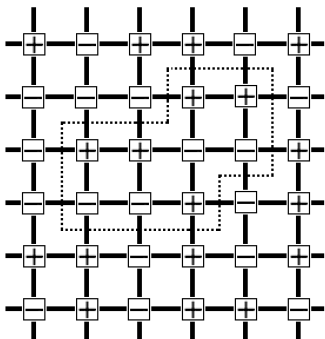
Let J be a configuration of couplings, sampled from the product measure ν .

- For any spin configuration α on \mathbb{Z}^2 , we say that α is an *infinite-volume ground state for J* if the energy of each closed dual loop is positive.
- The *energy* of a set S of dual edges in a coupling configuration J and a spin configuration α is the sum

$$\sum_{\langle x,y \rangle \in S} J_{xy} \alpha_x \alpha_y.$$

Definitions

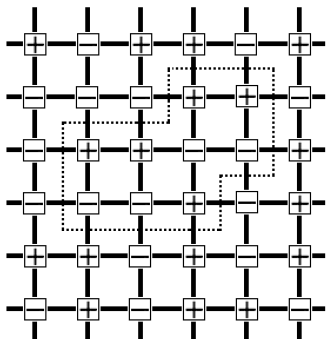
This coincides with the finite-volume definition of a ground state.



- In this case, the energy of a dual loop is $1/2$ of the energy needed to “flip” the spins surrounded by that loop.
- It costs positive energy to flip any finite set of spins.

Definitions

This coincides with the finite-volume definition of a ground state.



- Note that if α is a ground state then $-\alpha$ (the global flip of α) is also a ground state.
- Therefore we consider *ground state pairs* (GSP's).

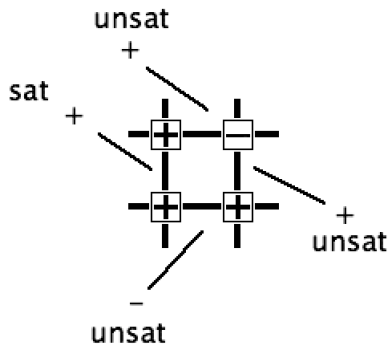
Domain walls

We begin with a metastate measure \mathcal{K}_J for \mathbb{Z}^2 (or for the half-plane).

- This is constructed using boxes $\Lambda^{(L)}$ (or Λ_n) of the form $[-n, n] \times [-n, n]$ with periodic boundary conditions (or periodic horizontally and free vertically).
- Recall that this is a measure on infinite-volume ground states (globally flip-related spin configurations) on \mathbb{Z}^2 (or the half-plane).

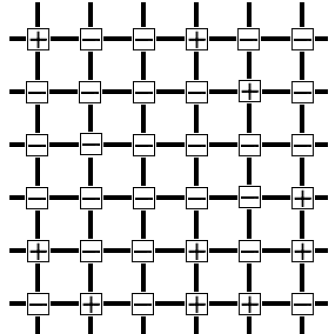
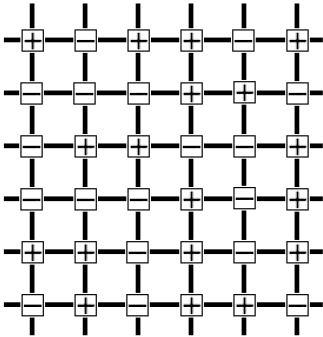
Domain walls

Sample two states (pairs) α and β from $\mathcal{K}_J^\#$. We define the *interface* $\alpha\Delta\beta$ as a subset of the dual lattice:



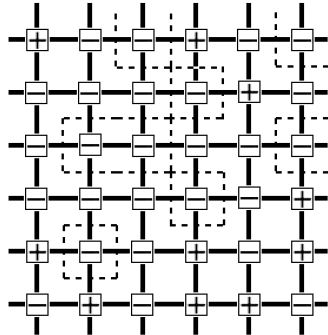
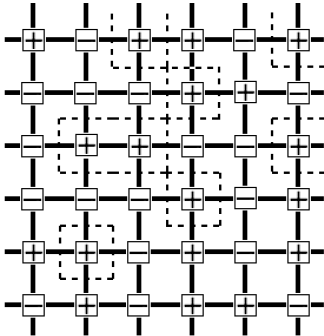
- a dual edge e^* is in the interface if it is satisfied in exactly one of α and β ; and
- e^* is *satisfied* in α if it has positive energy in α .

Example



- The interface separates regions where spins agree from where spins disagree.
- Therefore the connected components (*domain walls*) form either infinite paths or closed loops (no dangling ends).

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Domain walls for ground states

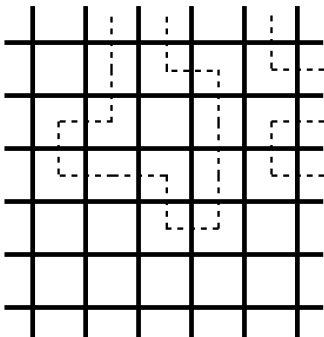
We want to study properties of domain walls between ground states. To do this:

- identify each state (spin config. pair) with its set of satisfied and unsatisfied couplings,
- choose two states from the product measure $\mathcal{K}_J \times \mathcal{K}_J$ and integrate out the couplings.

This gives a measure \mathcal{D} on domain wall configurations. It is translation-invariant because each \mathcal{K}_J was built with periodic boundary conditions.

Domain walls for ground states

If the \mathcal{K}_J were not supported on GSP's, we could sample something like this from \mathcal{D} (the configuration in the last example):



But interfaces between GSP's cannot contain closed loops (reasons soon). The config. at left is now admissible.

Newman-Stein Theorem

For the rest of the talk we will discuss the geometry of ground state domain walls.

Theorem (Newman-Stein)

Let α and β be two GSP's on \mathbb{Z}^2 (or on the half-plane).

1. The interface $\alpha\Delta\beta$ contains no loops or dangling ends.
2. If α and β are sampled from \mathcal{K}_J , (on \mathbb{Z}^2) then for a.e. J , $\alpha\Delta\beta$ contains no "branching points." In other words, the domain walls are double-infinite paths in which each (dual) vertex has degree two.

Here "dangling end" means a (dual) vertex of degree one in the interface.

Proof idea

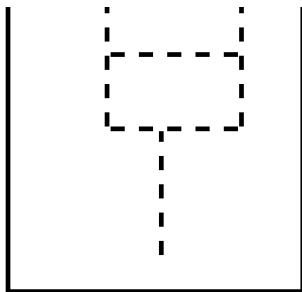
We saw that an interface between *any* two states cannot have dangling ends. (Interfaces form boundaries of sets of spins.) We must show:

- The interface between GSP's cannot have loops; and
- any interface sampled from \mathcal{D} (a.s.) cannot have branching points.

The first statement concerns *all* GSP interfaces. The second refers only to those constructed through the metastate (with p.b.c.)

Proof idea

In this figure, the interface has (a) a loop, (b) a dangling end, and (c) a branching point.

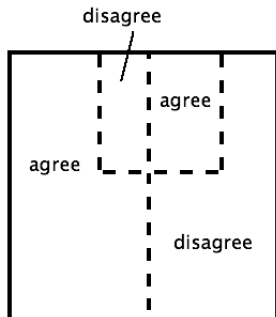


Loops

Branching points

Branching points

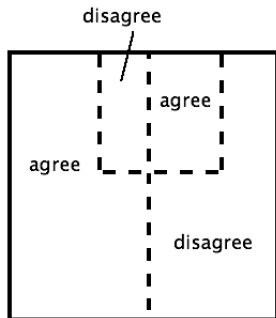
Now we must deal with branching points.



- Notice that since there are no loops, branching points must connect distinct paths to infinity.
- Also they must connect an even number (4) of paths.

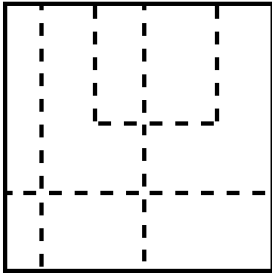
Branching points

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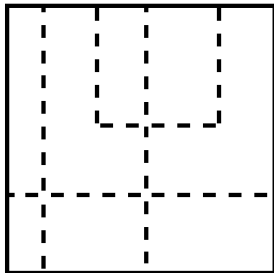
To rule out branching points
(dual vertices of degree ≥ 3)
we use the standard
Burton-Keane argument.

Branching points – Burton-Keane



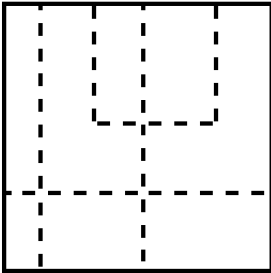
- We can associate with each branching point a unique point on the boundary. (Adding a branching point adds at least one new boundary point.)

Branching points – Burton-Keane



- This gives
number of branching points in box
 \leq surface area.

Branching points – Burton-Keane



- But by ergodicity, in expectation, number of branching points in box $\geq c$ (volume).
- This is a contradiction.

Final picture

Because of the theorem, if interfaces exist, they consist of disjoint doubly-infinite paths.