Extreme values in terms of a matrix exponential: splitting methods in the M/M/c retrial queue

M. López García Complutense University of Madrid, Spain

Talk based on the manuscript Maximum queue lengths during a fixed time interval in the M/M/c retrial queue, coauthored with A. Gómez-Corral

Organization of the talk

Introduction

Maximum queue length during a certain time interval

- Matrix exponential solution
- Computation of the matrix exponential solution
- Accuracy of the solution

3 Numerical results

4 Discussion

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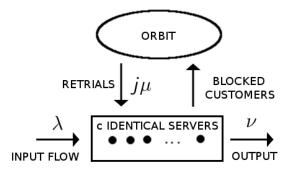
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The M/M/c retrial queue

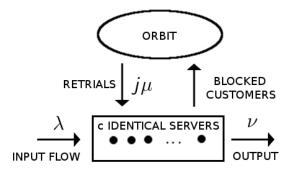
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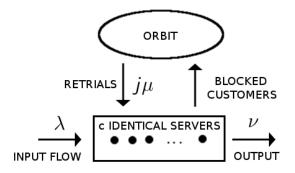
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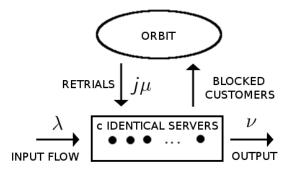
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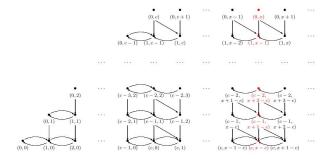
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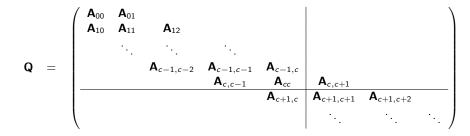
on $S = \{0, ..., c\} \times \{0, 1, 2, ...\}$, where C(t) is the number of busy servers, and N(t) is the number of customers in orbit at time t. If $S = \bigcup_{j=0}^{\infty} l(j)$ with $l(j) = \{(i, j - i) : 0 \le i \le \min(j, c)\}$, then \mathcal{X} is a LD-QBD process



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- Expected queue length.
- Waiting time.
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Assuming a stationary regime ($\rho = \lambda/(c\nu) < 1$) is needed to guarantee the existence, and subsequent computation, of most classical queueing measures.

Alternative: Maximum queue length distribution (X_{max}) in a busy period (the period that starts when the process leaves the state (0,0) and ends at the first epoch thereafter that the process visits the state (0,0) again).

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- It is a measure of system congestion.
- It gives support to the adoption of drastic decisions such as an increase of the number of agents or rescheduling of common resources.
- It can be computed even in non-stationary regime, but it might be defective (i.e., P(X_{max} < ∞) < 1).

- Artalejo JR, Economou A, López Herrero MJ. "Algorithmic analysis of the maximum queue length in a busy period for the M/M/c retrial queue". INFORMS Journal on Computing, in revision.
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Objective Computation of $P(Z(t_0) \le x | C(0) = i, N(0) = j), \text{ for } x \ge i + j \text{ and } (i, j) \in S,$ since $P(Z(t_0) \le x | C(0) = i, N(0) = j) = 0 \text{ if } x < i + j.$ Instead of X_{max} , we may be interested in

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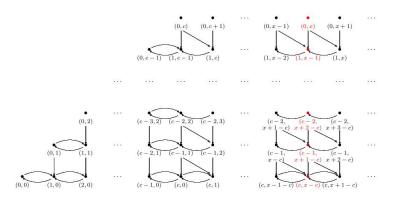
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Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

For $x \ge i + j$ and $(i, j) \in S$, the conditional probability $P(Z(t_0) > x | C(0) = i, N(0) = j)$

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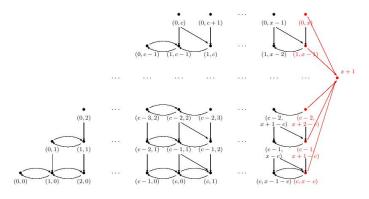
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We consider an auxiliary absorbing process $\overline{\mathcal{X}}(x) = \{\overline{X}(t) : t \ge 0\}$ defined on the state space

$$\overline{\mathcal{S}}(x) = \bigcup_{k=0}^{x} l(k) \cup \{x+1\},$$

where the kth level I(k) is given by

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$$\overline{\mathbf{Q}}(x) = \begin{pmatrix} \mathbf{A}_{00} & \mathbf{A}_{01} & & & \\ \mathbf{A}_{10} & \mathbf{A}_{11} & \mathbf{A}_{12} & & & \\ & \ddots & \ddots & \ddots & \\ & \mathbf{A}_{x-1,x-2} & \mathbf{A}_{x-1,x-1} & \mathbf{A}_{x-1,x} \\ & & \mathbf{A}_{x,x-1} & \mathbf{A}_{xx} & \lambda \mathbf{e}_{\#I(x)} \\ \hline & & & 0 \end{pmatrix}$$
$$= \left(\frac{\mathbf{T}(x) \mid \mathbf{t}_{x+1}(x)}{\mathbf{0}_{J(x)}^{T} \mid 0} \right), \quad x \leq c, \text{ or } x \geq c+1?$$

and its standard transition function $\overline{\mathbf{P}}(t_0; x)$ can be expressed as

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Therefore, we have

$$P(Z(t_0) \le x | C(0) = i, N(0) = j) = 1 - \overline{e}_{J(x)}(i, j) \mathbf{p}_{x+1}(t_0; x),$$

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$$P(Z(t_0) \le x | C(0) = i, N(0) = j) = 1 - \overline{e}_{J(x)}(i, j) \left(I_{J(x)} - \exp\{T(x)t_0\} \right) \left(-T^{-1}(x) \right) t_{x+1}(x).$$

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Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

In principle, $\exp{\{\mathbf{T}(x)t_0\}}$ could be computed in many ways:

- Series methods: Taylor series, Padé approximation, scaling and squaring, Chebyshev rational approximation
- Ordinary differential equation methods: general purpose O.D.E. solver, single step / multistep O.D.E. solvers
- Polynomial methods: Cayley-Hamilton, Lagrange interpolation, Newton interpolation, Vandermonde, inverse Laplace transforms, companion matrix
- Matrix decomposition methods: eigenvectors, triangular systems of eigenvectors, Jordan canonical form, Schur, block diagonal
- Splitting methods

An interesting survey is the paper by Moler and Van Loan (2003) Software: MathLab, Mathematica, Mapple, ISML library, etc.

In practice, consideration of computational stability and efficiency indicates that some of the methods are preferable to others, but that none are completely satisfactory when they are implemented as general-purpose algorithms

For a value $x \ge i + j$, the dimension of $\exp\{\mathbf{T}(x)t_0\}$ is given by $J(x) = \frac{(x+1)(x+2)}{2}$ if $1 \le x \le c$, and $\frac{(c+1)(c+2)}{2} + (x-c)(c+1)$ if $x \ge c+1$.

As a result,

- Increasing values of x will imply more demanding memory requirements
- General-purpose algorithms will fail to give satisfactory results as *x* progressively increases

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

Splitting methods

For a certain splitting $\mathbf{T}(x) = \mathbf{U}(x) + \mathbf{V}(x)$, it is known that

 $\exp\{\mathbf{T}(x)t_0\} = \exp\{\mathbf{U}(x)t_0\}\exp\{\mathbf{V}(x)t_0\}$

if and only if $\mathbf{U}(x)$ and $\mathbf{V}(x)$ commute

As $\mathbf{U}(x)$ and $\mathbf{V}(x)$ do not commute, the exponentials of the matrices $\mathbf{U}(x)$ and $\mathbf{V}(x)$ are directly related to that of $\mathbf{T}(x)$ by

$$\exp\{\mathbf{T}(x)t_0\} = \lim_{p \to \infty} \left(\exp\left\{\mathbf{U}(x)\frac{t_0}{p}\right\}\exp\left\{\mathbf{V}(x)\frac{t_0}{p}\right\}\right)^p.$$

Moller and Van Loan (2003) suggest the approximation

 $\exp\{\mathbf{T}(x)t_0\} \simeq (\exp\{\mathbf{U}(x)t\}\exp\{\mathbf{V}(x)t\})^{p_0},$

where $t = p_0^{-1} t_0$, for an appropriately selected integer p_0 .

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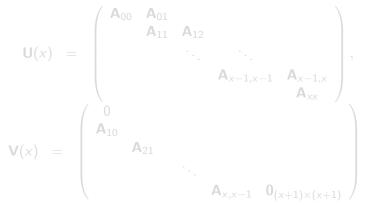
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$exp\{\mathbf{T}(x)t_0\}$

Case 1: $1 \le x \le c$. T(x) = U(x) + V(x) is defined by

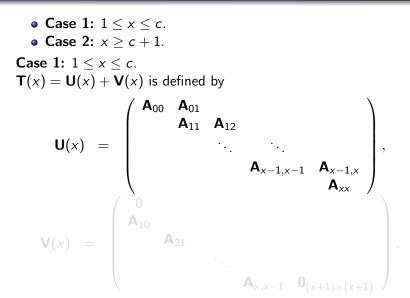


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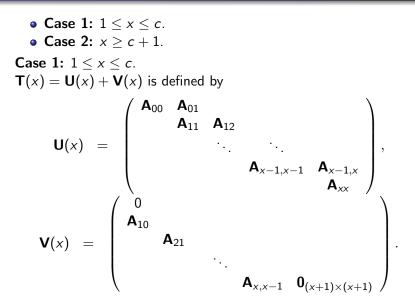
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Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

$exp\{\mathbf{T}(x)t_0\}$



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For $1 \le x \le c$, it is seen that

$$\exp\{\mathbf{V}(x)t\} = \sum_{k=0}^{x} \frac{t^{k}}{k!} \mathbf{V}^{k}(x)$$
$$= \begin{pmatrix} \mathbf{M}(0;0) & \\ \mathbf{M}(1;0) & \mathbf{M}(1;1) & \\ \vdots & \vdots & \ddots & \\ \mathbf{M}(x-1;0) & \mathbf{M}(x-1;1) & \cdots & \mathbf{M}(x-1;x-1) & \\ \mathbf{M}(x;0) & \mathbf{M}(x;1) & \cdots & \mathbf{M}(x;x-1) & \mathbf{M}(x;x) \end{pmatrix}$$

where

$$\mathsf{M}(y;y') = \begin{cases} \mathsf{I}_{y+1} & \text{if } 1 \le y \le x, \\ \mathsf{A}_{y,y-1} \cdot \ldots \cdot \mathsf{A}_{y'+1,y'}, & \text{if } 0 \le y' \le y-1, 1 \le y \le x. \end{cases}$$

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For $1 \le x \le c$, it is seen that

$$\exp\{\mathbf{V}(x)t\} = \sum_{k=0}^{x} \frac{t^{k}}{k!} \mathbf{V}^{k}(x)$$
$$= \begin{pmatrix} \mathbf{M}(0;0) & \\ \mathbf{M}(1;0) & \mathbf{M}(1;1) & \\ \vdots & \vdots & \ddots & \\ \mathbf{M}(x-1;0) & \mathbf{M}(x-1;1) & \cdots & \mathbf{M}(x-1;x-1) & \\ \mathbf{M}(x;0) & \mathbf{M}(x;1) & \cdots & \mathbf{M}(x;x-1) & \mathbf{M}(x;x) \end{pmatrix}$$

where

$$\mathsf{M}(y;y') = \begin{cases} \mathsf{I}_{y+1} & \text{if } 1 \leq y \leq x, \\ \mathsf{A}_{y,y-1} \cdot \ldots \cdot \mathsf{A}_{y'+1,y'}, & \text{if } 0 \leq y' \leq y-1, 1 \leq y \leq x. \end{cases}$$

,

Technical condition: For $1 \le x \le c$, the eigenvalues of the matrix $\mathbf{U}(x)$ are all distinct if and only if

• (A.1) [Eigenvalues of the sub-matrix **A**_{yy} are all distinct]:

$$\nu \neq \mu$$

(A.2) [Eigenvalues of sub-matrices A_{yy} and A_{y'y'} with y < y' are also distinct]:

$$\nu \neq \left(1 + \frac{y' - y}{I - I'}\right)\mu,$$

for every pair (y, y') of integers with $0 \le y < y' \le x$, and integers $0 \le l \le y$ and $0 \le l' \le y'$ with $l \ne l'$

Under the technical condition, we derive the decomposition formula $\mathbf{U}(x) = \mathbf{R}_x \operatorname{diag}[r(0; 0), ..., r(x; x)]\mathbf{R}_x^{-1}$, which implies

$$\exp\{\mathbf{U}(x)t\} = \mathbf{R}_{x} \begin{pmatrix} e^{r(0;0)t} & & \\ & e^{r(1;0)t} & & \\ & & e^{r(1;1)t} & \\ & & & \ddots \\ & & & & e^{r(x;x)t} \end{pmatrix} \mathbf{R}_{x}^{-1},$$

where $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$ and \mathbf{R}_x consists of the right eigenvectors of $\mathbf{U}(x)$ associated with the eigenvalues r(y; l), for $0 \le l \le y \le x$.

The matrix \mathbf{R}_{x} has the structured form

$$\mathbf{R}_{x} \; = \; \begin{pmatrix} \begin{array}{c|cccccc} \mathbf{P}_{00} & \mathbf{P}_{01} & \mathbf{P}_{02} & \cdots & \mathbf{P}_{0,x-1} & \mathbf{P}_{0x} \\ & \mathbf{P}_{11} & \mathbf{P}_{12} & \cdots & \mathbf{P}_{1,x-1} & \mathbf{P}_{1x} \\ & \mathbf{P}_{22} & \cdots & \mathbf{P}_{2,x-1} & \mathbf{P}_{2x} \\ & & \ddots & \vdots & \vdots \\ & & & \mathbf{P}_{x-1,x-1} & \mathbf{P}_{x-1,x} \\ \hline & & & & & \mathbf{P}_{xx} \end{pmatrix} \; = \; \left(\begin{array}{c|c} \mathbf{R}_{x-1} & \mathbf{N}_{x} \\ \hline & & \mathbf{P}_{xx} \end{array} \right),$$

where the columns in $\mathbf{P}_{y'y} = [\mathbf{p}(y', y; 0), \mathbf{p}(y', y; 1), ..., \mathbf{p}(y', y; y)]$ are given by

$$\mathbf{p}(y',y;l) = \begin{cases} \mathbf{v}(y;l), & \text{if } y' = y, \\ \prod_{k=y'}^{y-1} (r(y;l)\mathbf{I}_{k+1} - \mathbf{A}_{kk})^{-1} \mathbf{A}_{k,k+1} \mathbf{v}(y;l), & \text{if } 0 \le y' \le y - 1. \end{cases}$$

and $\mathbf{v}(y; l)$ denotes the right eigenvector of \mathbf{A}_{yy} associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$.

$$(r(y;l)\mathbf{I}_{k+1} - \mathbf{A}_{kk})^{-1} = \sum_{l'=0}^{k} \frac{\mathbf{v}(k;l')\mathbf{w}(k;l')}{r(y;l) - r(k;l')}, \quad 1 \le k \le c$$

The matrix \mathbf{R}_{x} has the structured form

$$\mathbf{R}_{x} \; = \; \begin{pmatrix} \begin{array}{c|cccccc} \mathbf{P}_{00} & \mathbf{P}_{01} & \mathbf{P}_{02} & \cdots & \mathbf{P}_{0,x-1} & \mathbf{P}_{0x} \\ & \mathbf{P}_{11} & \mathbf{P}_{12} & \cdots & \mathbf{P}_{1,x-1} & \mathbf{P}_{1x} \\ & \mathbf{P}_{22} & \cdots & \mathbf{P}_{2,x-1} & \mathbf{P}_{2x} \\ & & \ddots & \vdots & \vdots \\ & & & \mathbf{P}_{x-1,x-1} & \mathbf{P}_{x-1,x} \\ \hline & & & & & \mathbf{P}_{xx} \end{pmatrix} \; = \; \left(\begin{array}{c|c} \mathbf{R}_{x-1} & \mathbf{N}_{x} \\ \hline & & \mathbf{P}_{xx} \end{array} \right),$$

where the columns in $\mathbf{P}_{y'y} = [\mathbf{p}(y', y; 0), \mathbf{p}(y', y; 1), ..., \mathbf{p}(y', y; y)]$ are given by

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and $\mathbf{v}(y; l)$ denotes the right eigenvector of \mathbf{A}_{yy} associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$.

$$(r(y;l)\mathbf{I}_{k+1} - \mathbf{A}_{kk})^{-1} = \sum_{l'=0}^{k} \frac{\mathbf{v}(k;l')\mathbf{w}(k;l')}{r(y;l) - r(k;l')}, \quad 1 \le k \le c$$

Right eigenvector $\mathbf{v}(y; l)$ of \mathbf{A}_{yy} with $0 \le y \le c$ associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$, for $0 \le l \le y$:

$$v_{l'}(y;l) = \begin{cases} \binom{y-l'}{y-l} \left(1 - \frac{\nu}{\mu}\right)^{-(l-l')}, & \text{if } 0 \le l' \le l-1, \\ 1, & \text{if } l' = l, \\ 0, & \text{if } l+1 \le l' \le y. \end{cases}$$

Left eigenvector $\mathbf{w}(y; l)$ of \mathbf{A}_{yy} with $0 \le y \le c$, associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$, for $0 \le l \le y$:

$$w_{l'}(y;l) = \left\{egin{array}{ll} 0, & ext{if } 0 \leq l' \leq l-1, \ 1, & ext{if } l' = l, \ \left(rac{y-l}{y-l'}
ight) \left(rac{
u}{\mu} - 1
ight)^{-(l'-l)}, & ext{if } l+1 \leq l' \leq y. \end{array}
ight.$$

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

Starting with $\exp{\{\mathbf{U}(0)t\}} = e^{-\lambda t}$, the matrix exponential $\exp{\{\mathbf{U}(x)t\}}$ is evaluated from

$$\exp\{\mathbf{U}(x)t\} = \begin{pmatrix} \exp\{\mathbf{U}(x-1)t\} & \mathbf{N}_{x}(t) \\ \mathbf{0}_{(x+1)\times J(x-1)} & \exp\{\mathbf{A}_{xx}t\} \end{pmatrix}, \quad 1 \le x \le c,$$

where $\exp{\{\mathbf{A}_{xx}t\}} = \mathbf{P}_{xx}\mathbf{T}_{xx}(t)\mathbf{P}_{xx}^{-1}$ with

$$\begin{split} \mathbf{N}_{x}(t) &= (\mathbf{N}_{x}\mathbf{T}_{xx}(t) - \exp\{\mathbf{U}(x-1)t\}\mathbf{N}_{x}) \, \mathbf{P}_{xx}^{-1} \\ \mathbf{T}_{xx}(t) &= \text{diag}[e^{r(x;0)t}, e^{r(x;1)t}, ..., e^{r(x;x)t}], \\ \mathbf{P}_{xx} &= [\mathbf{v}(x;0), \mathbf{v}(x;1), ..., \mathbf{v}(x;x)], \\ \mathbf{P}_{xx}^{-1} &= \begin{pmatrix} c^{-1}(x;0)\mathbf{w}(x;0) \\ c^{-1}(x;1)\mathbf{w}(x;1) \\ \vdots \\ c^{-1}(x;x)\mathbf{w}(x;x) \end{pmatrix}, \end{split}$$

and $c(x; l) = \sum_{k=0}^{x} (\mathbf{w}(x; l))_k (\mathbf{v}(x; l))_k$.

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

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$$\exp\{\mathbf{U}(x)t\} = \begin{pmatrix} \exp\{\mathbf{U}(x-1)t\} & \mathbf{N}_{x}(t) \\ \mathbf{0}_{(x+1)\times J(x-1)} & \exp\{\mathbf{A}_{xx}t\} \end{pmatrix}, \quad 1 \le x \le c,$$

where $\exp{\{\mathbf{A}_{xx}t\}} = \mathbf{P}_{xx}\mathbf{T}_{xx}(t)\mathbf{P}_{xx}^{-1}$ with

$$\begin{split} \mathbf{N}_{x}(t) &= (\mathbf{N}_{x}\mathbf{T}_{xx}(t) - \exp\{\mathbf{U}(x-1)t\}\mathbf{N}_{x})\mathbf{P}_{xx}^{-1}, \\ \mathbf{T}_{xx}(t) &= \operatorname{diag}[e^{r(x;0)t}, e^{r(x;1)t}, ..., e^{r(x;x)t}], \\ \mathbf{P}_{xx} &= [\mathbf{v}(x;0), \mathbf{v}(x;1), ..., \mathbf{v}(x;x)], \\ \mathbf{P}_{xx}^{-1} &= \begin{pmatrix} c^{-1}(x;0)\mathbf{w}(x;0) \\ c^{-1}(x;1)\mathbf{w}(x;1) \\ \vdots \\ c^{-1}(x;x)\mathbf{w}(x;x) \end{pmatrix}, \end{split}$$

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Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

Starting with $\exp{\{\mathbf{U}(0)t\}} = e^{-\lambda t}$, the matrix exponential $\exp{\{\mathbf{U}(x)t\}}$ is evaluated from

$$\exp\{\mathbf{U}(x)t\} = \begin{pmatrix} \exp\{\mathbf{U}(x-1)t\} & \mathbf{N}_{x}(t) \\ \mathbf{0}_{(x+1)\times J(x-1)} & \exp\{\mathbf{A}_{xx}t\} \end{pmatrix}, \quad 1 \le x \le c,$$

where $\exp{\{\mathbf{A}_{xx}t\}} = \mathbf{P}_{xx}\mathbf{T}_{xx}(t)\mathbf{P}_{xx}^{-1}$ with

$$\begin{split} \mathbf{N}_{x}(t) &= (\mathbf{N}_{x}\mathbf{T}_{xx}(t) - \exp\{\mathbf{U}(x-1)t\}\mathbf{N}_{x})\mathbf{P}_{xx}^{-1}, \\ \mathbf{T}_{xx}(t) &= \operatorname{diag}[e^{r(x;0)t}, e^{r(x;1)t}, ..., e^{r(x;x)t}], \\ \mathbf{P}_{xx} &= [\mathbf{v}(x;0), \mathbf{v}(x;1), ..., \mathbf{v}(x;x)], \\ \mathbf{P}_{xx}^{-1} &= \begin{pmatrix} c^{-1}(x;0)\mathbf{w}(x;0) \\ c^{-1}(x;1)\mathbf{w}(x;1) \\ \vdots \\ c^{-1}(x;x)\mathbf{w}(x;x) \end{pmatrix}, \end{split}$$

and $c(x; l) = \sum_{k=0}^{x} (\mathbf{w}(x; l))_{k} (\mathbf{v}(x; l))_{k}$.

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

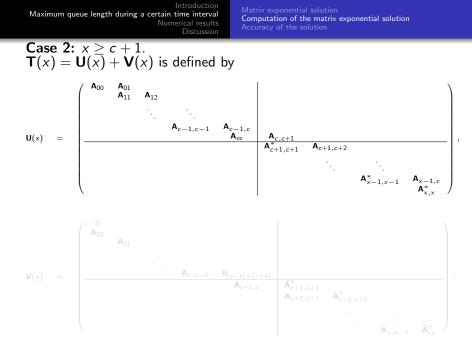
Starting with $\exp{\{\mathbf{U}(0)t\}} = e^{-\lambda t}$, the matrix exponential $\exp{\{\mathbf{U}(x)t\}}$ is evaluated from

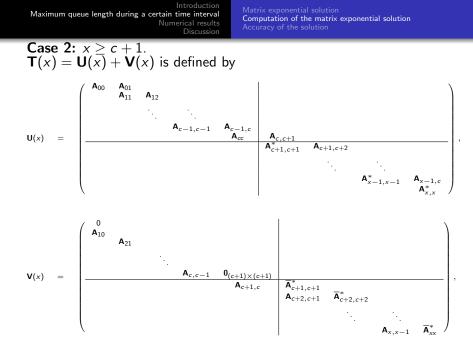
$$\exp\{\mathbf{U}(x)t\} = \begin{pmatrix} \exp\{\mathbf{U}(x-1)t\} & \mathbf{N}_{x}(t) \\ \mathbf{0}_{(x+1)\times J(x-1)} & \exp\{\mathbf{A}_{xx}t\} \end{pmatrix}, \quad 1 \le x \le c,$$

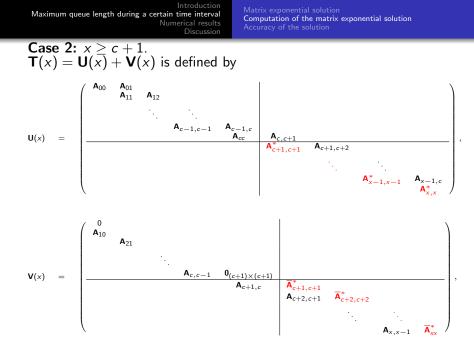
where $\exp{\{\mathbf{A}_{xx}t\}} = \mathbf{P}_{xx}\mathbf{T}_{xx}(t)\mathbf{P}_{xx}^{-1}$ with

$$\begin{split} \mathbf{N}_{x}(t) &= (\mathbf{N}_{x}\mathbf{T}_{xx}(t) - \exp\{\mathbf{U}(x-1)t\}\mathbf{N}_{x})\mathbf{P}_{xx}^{-1}, \\ \mathbf{T}_{xx}(t) &= \operatorname{diag}[e^{r(x;0)t}, e^{r(x;1)t}, ..., e^{r(x;x)t}], \\ \mathbf{P}_{xx} &= [\mathbf{v}(x;0), \mathbf{v}(x;1), ..., \mathbf{v}(x;x)], \\ \mathbf{P}_{xx}^{-1} &= \begin{pmatrix} c^{-1}(x;0)\mathbf{w}(x;0) \\ c^{-1}(x;1)\mathbf{w}(x;1) \\ \vdots \\ c^{-1}(x;x)\mathbf{w}(x;x) \end{pmatrix}, \end{split}$$

and $c(x; l) = \sum_{k=0}^{x} (\mathbf{w}(x; l))_{k} (\mathbf{v}(x; l))_{k}$.







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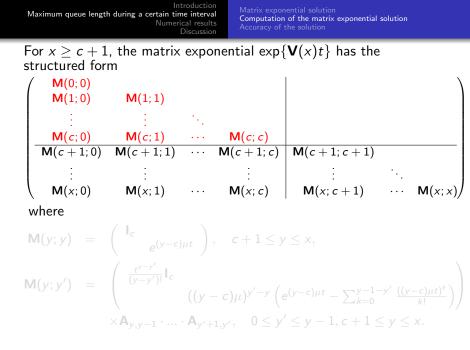
where

$$\overline{\mathbf{A}}_{yy}^{*} = \begin{pmatrix} 0 & 0 & \cdots & 0 & 0 \\ 0 & 0 & \cdots & 0 & 0 \\ \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & \cdots & 0 & 0 \\ 0 & 0 & \cdots & 0 & (y-c)\mu \end{pmatrix}, \quad y \ge c+1,$$

and the matrix \boldsymbol{A}_{yy}^{*} is given by

$$\begin{pmatrix} -\lambda - y\mu & y\mu \\ & -\lambda - \nu - (y-1)\mu & (y-1)\mu \\ & \ddots & \ddots \\ & & \ddots \\ & & -\lambda - (c-1)\nu - (y-c+1)\mu & (y-c+1)\mu \\ & & -\lambda - c\nu - (y-c)\mu \end{pmatrix}$$

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$$x \ge c + 1$$
, the matrix exponential exp $\{V(x)t\}$ has the
structured form $\left(\begin{array}{c|c} M(0;0) \\ M(1;0) & M(1;1) \\ \vdots & \vdots & \ddots \\ M(c;0) & M(c;1) & \cdots & M(c;c) \\ \hline M(c+1;0) & M(c+1;1) & \cdots & M(c+1;c) \\ \vdots & \vdots & \ddots \\ M(x;0) & M(x;1) & \cdots & M(x;c) \\ \hline M(x;c+1) & \cdots & M(x;x) \\ \end{array} \right)$ where $M(y;y) = \left(\begin{array}{c|c} I_c \\ e^{(y-c)\mu t} \\ (y'-y') \end{array} \right), \quad c+1 \le y \le x,$ $M(y;y') = \left(\begin{array}{c|c} I_c \\ e^{(y-c)\mu t} \\ (y'-y') \end{array} \right), \quad c+1 \le y \le x,$ $M(y;y') = \left(\begin{array}{c|c} I_c \\ e^{(y-c)\mu t} \\ (y'-y') \end{array} \right), \quad c < 1 \le y \le x,$ $M(y;y') = \left(\begin{array}{c|c} I_c \\ e^{(y-c)\mu t} \\ (y'-y') \end{array} \right), \quad c < 1 \le y \le x,$ $M(y;y') = \left(\begin{array}{c|c} I_c \\ e^{(y-c)\mu t} \\ (y'-y') \end{array} \right), \quad c < 1 \le y \le x.$



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structured formFor $x \ge c + 1$, the matrix exponential exp $\{V(x)t\}$ has the
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 $\times \mathbf{A}_{y,y-1} \cdot \ldots \cdot \mathbf{A}_{y'+1,y'}, \quad 0 \leq y' \leq y-1, c+1 \leq y \leq x.$

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, the matrix exponential exp $\{V(x)t\}$ has the
structured formFor $x \ge c + 1$, the matrix exponential exp $\{V(x)t\}$ has the
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For values $x \ge c + 1$,

$$\exp\{\mathbf{U}(x)t\} = \begin{pmatrix} \exp\{\mathbf{U}(x-1)t\} & \mathbf{N}_{x}(t) \\ \mathbf{0}_{(c+1)\times J(x-1)} & \exp\{\mathbf{A}_{xx}^{*}t\} \end{pmatrix}, \quad x \ge c,$$

where $\exp{\{\mathbf{A}_{xx}^{*}t\}} = \mathbf{P}_{xx}\mathbf{T}_{xx}(t)\mathbf{P}_{xx}^{-1}$ with

$$\begin{split} \mathbf{N}_{x}(t) &= (\mathbf{N}_{x}\mathbf{T}_{xx}(t) - \exp\{\mathbf{U}(x-1)t\}\mathbf{N}_{x})\mathbf{P}_{xx}^{-1} \\ \mathbf{T}_{xx}(t) &= \operatorname{diag}[e^{r(x;0)t}, e^{r(x;1)t}, ..., e^{r(x;c)t}], \\ \mathbf{P}_{xx} &= [\mathbf{v}(x;0), \mathbf{v}(x;1), ..., \mathbf{v}(x;c)], \\ \mathbf{P}_{xx}^{-1} &= \begin{pmatrix} c^{-1}(x;0)\mathbf{w}(x;0) \\ c^{-1}(x;1)\mathbf{w}(x;1) \\ \vdots \\ c^{-1}(x;x)\mathbf{w}(x;c) \end{pmatrix}, \end{split}$$

and $c(x; l) = \sum_{k=0}^{c} (\mathbf{w}(x; l))_{k} (\mathbf{v}(x; l))_{k}$.

Right eigenvector $\mathbf{v}(y; l)$ of \mathbf{A}_{yy} with $0 \le y \le c$ and \mathbf{A}^*_{yy} with $c + 1 \le y$, associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$, for $0 \le l \le \min(y, c)$:

$$v_{l'}(y; l) = \begin{cases} \binom{y-l'}{y-l} \left(1 - \frac{\nu}{\mu}\right)^{-(l-l')}, & \text{if } 0 \le l' \le l-1, \\ 1, & \text{if } l' = l, \\ 0, & \text{if } l+1 \le l' \le \min(y, c). \end{cases}$$

Left eigenvector $\mathbf{w}(y; l)$ of \mathbf{A}_{yy} with $0 \le y \le c$ and \mathbf{A}_{yy}^* with $c + 1 \le y$, associated with the eigenvalue $r(y; l) = -(\lambda + l\nu + (y - l)\mu)$, for $0 \le l \le \min(y, c)$: $w_{l'}(y; l) = \begin{cases} 0, & \text{if } 0 \le l' \le l - 1, \\ 1, & \text{if } l' = l, \\ \binom{y-l}{y-l'} \left(\frac{\nu}{\mu} - 1\right)^{-(l'-l)}, & \text{if } l + 1 \le l' \le \min(y, c). \end{cases}$ For values $x \ge 1$, the technical condition becomes

The eigenvalues of $\mathbf{U}(x)$ are all distinct if and only if $\nu \neq \mu$ and

$$u \neq \left(1 + \frac{y' - y}{I - I'}\right)\mu,$$

for every pair of integers $0 \le y < y' \le x$, and every integers $0 \le l \le \min(y, c)$ and $0 \le l' \le \min(y', c)$ with $l \ne l'$.

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

Selection of p_0 : In the approximation

$$\exp\{\mathbf{T}(x)t_0\} \quad \simeq \quad \left(\exp\{\mathbf{U}(x)t\}\exp\{\mathbf{V}(x)t\}\right)^{p_0},$$

where $t = p_0^{-1} t_0$, we choose p_0 such that

$$\left|\left|\exp\{\mathbf{T}(x)t_0\}-(\exp{\{\mathbf{U}(x)t\}}\exp{\{\mathbf{V}(x)t\}})^{p_0}\right|\right|_{\infty} < \varepsilon,$$

where the maximum row sum matrix norm is defined by $||\mathbf{W}||_{\infty} = \max_{1 \le i \le k} \sum_{j=1}^{k} |w_{ij}|.$ Then, it is seen that

$$\left|\left|\exp\{\mathsf{T}(x)t_0\} - \left(\exp\{\mathsf{U}(x)t\}\exp\{\mathsf{V}(x)t\}\right)^{p_0}\right|\right|_{\infty} \le \frac{h(x;t_0)}{2p_0},$$

where $\mu_{\infty}(\cdot)$ is the logarithmic norm of a matrix and

 $h(x; t_0) = ||\mathbf{U}(x)\mathbf{V}(x) - \mathbf{V}(x)\mathbf{U}(x)||_{\infty} t_0^2 e^{(\mu_{\infty}(\mathbf{U}(x)) + \mu_{\infty}(\mathbf{V}(x)))t_0}.$

Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

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Matrix exponential solution Computation of the matrix exponential solution Accuracy of the solution

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$$\left|\left|\exp\{\mathbf{T}(x)t_0\}-\left(\exp\{\mathbf{U}(x)t\}\exp\{\mathbf{V}(x)t\}\right)^{p_0}\right|\right|_{\infty} \le \frac{h(x;t_0)}{2p_0},$$

where $\mu_{\infty}(\cdot)$ is the logarithmic norm of a matrix and

$$h(x; t_0) = ||\mathbf{U}(x)\mathbf{V}(x) - \mathbf{V}(x)\mathbf{U}(x)||_{\infty} t_0^2 e^{(\mu_{\infty}(\mathbf{U}(x)) + \mu_{\infty}(\mathbf{V}(x)))t_0}.$$

In particular,

$$\begin{array}{lll} \mu_{\infty}(\mathbf{U}(x)) &=& 0, \quad x \geq 1, \\ \mu_{\infty}(\mathbf{V}(x)) &=& \left\{ \begin{array}{ll} x\nu, & \text{if } 1 \leq x \leq c, \\ c\nu + (x-c)\mu, & \text{if } c+1 \leq x, \end{array} \right. \end{array}$$

For $1 \le x \le c$, it is readily seen that

$$||\mathbf{U}(x)\mathbf{V}(x) - \mathbf{V}(x)\mathbf{U}(x)||_{\infty} = \begin{cases} \nu\mu x, & \text{if } \lambda + \nu \leq \mu, \\ \nu(\lambda + \nu)x, & \text{if } \lambda + \nu > \mu. \end{cases}$$

(Similar expressions for $||\mathbf{U}(x)\mathbf{V}(x) - \mathbf{V}(x)\mathbf{U}(x)||_{\infty}$ in the case $c + 1 \le x$ are readily obtained.)

In the expression for $P(Z(t_0) \le x | C(0) = i, N(0) = j)$, that is,

$$1 - \bar{\mathbf{e}}_{J(x)}(i,j) \left(\mathbf{I}_{J(x)} - \exp\{\mathbf{T}(x)t_0\} \right) \mathbf{e}_{J(x)},$$

we suggest to replace $\exp{\{\mathbf{T}(x)t_0\}}$ by

 $\left(\exp\{\mathbf{U}(x)t\}\exp\{\mathbf{V}(x)t\}\right)^{p_0},$

$$\begin{aligned} \left| P(Z(t_0) \le x | C(0) = i, N(0) = j) - \hat{P}(Z(t_0) \le x | C(0) = i, N(0) = j) \right| \\ \le \left| \left| \exp\{\mathbf{T}(x)t_0\} - (\exp\{\mathbf{U}(x)t\} \exp\{\mathbf{V}(x)t\})^{p_0} \right| \right|_{\infty} \\ \le \frac{h(x; t_0)}{2p_0} < \varepsilon. \end{aligned}$$

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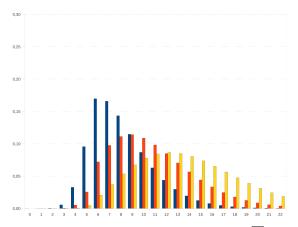
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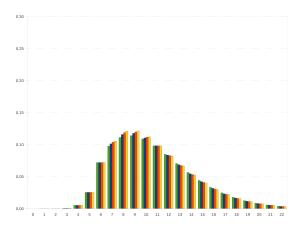
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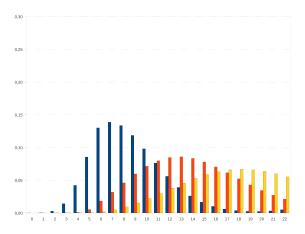
$$\begin{split} & \left| P(Z(t_0) \le x | C(0) = i, N(0) = j) - \hat{P}(Z(t_0) \le x | C(0) = i, N(0) = j) \right| \\ & \le \left| \left| \exp\{\mathbf{T}(x)t_0\} - (\exp\{\mathbf{U}(x)t\} \exp\{\mathbf{V}(x)t\})^{p_0} \right| \right|_{\infty} \\ & \le \frac{h(x; t_0)}{2p_0} < \varepsilon. \end{split}$$



The mass function of $Z(t_0)$ versus ρ for retrial queues with c = 6, $\nu = 3.0\sqrt{2.0}$ and $\mu = 2.5$. Values of the traffic load: $\rho = 0.8$, 1.0 and 1.2 (from left to right); interval length: $t_0 = 1.0$; initial state: (i, j) = (0, 0)



The mass function of $Z(t_0)$ versus μ for retrial queues with c = 6, $\nu = 3.0\sqrt{2.0}$ and $\rho = 1.0$. Values of the retrial rate: $\mu = 2.5$, 5.0, 7.5 and 10.0 (from left to right); interval length: $t_0 = 1.0$; initial state: (i, j) = (0, 0)



The mass function of $Z(t_0)$ versus t_0 for retrial queues with c = 6, $\nu = 3.0\sqrt{2.0}$ and $\mu = 2.5$. Traffic load: $\rho = 1.5$; values of the interval length: $t_0 = 1/3$, 2/3 and 1.0 (from left to right); interval length: $t_0 = 1.0$; initial state: (i, j) = (0, 0)

- It has always got a non-defective distribution, even if the LD-QBD process is not positive recurrent.
- Its probability distribution function has a matrix exponential form.
- We present simple conditions on the service rate ν and the retrial rate μ for the matrix exponential solution to be explicit or algorithmically tractable.
- We present an iterative scheme for computing the matrix exponential solution.
- A particularly appealing feature of this iterative solution based on splitting methods and eigenvalues/eigenvectors is that it allows us to obtain global error control.

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Thank you for your attention

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