A comparison between different cycle decompositions for Metropolis dynamics

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Abstract. In the last decades the problem of metastability has been attacked on rigorous grounds via many different approaches and techniques which are briefly reviewed in this paper. It is then useful to understand connections between different point of views. In view of this we consider irreducible, aperiodic and reversible Markov chains with exponentially small transition probabilities in the framework of Metropolis dynamics. We compare two different cycle decompositions and prove their equivalence.

Keywords: Stochastic dynamics, Markov chains, hitting times, metastability, Metropolis dynamics

MSC2010: 60K35,82C26

1. Introduction

Cycle decomposition is a very useful tool to estimate first hitting times for stochastic processes. In this note we compare, and prove the equivalence, between two different approaches that will be respectively called *path* and *graph cycle* decompositions. These results are utterly important in the generic study of the metastability phenomenon.

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1.1. Metastability

The phenomenon of metastability is defined by the following scenario: (i) a system is “trapped” for an abnormally long time in a state — the metastable state — different from the eventual equilibrium state consistent with the thermodynamical parameters. (ii) In the metastable state the system behaves as if it were in regular equilibrium. (iii) Subsequently, the system undergoes a sudden transition at a random time from the metastable to the stable state. The mathematical study of this phenomenon has been a standing issue since the foundation of Statistical Mechanics, but only in the 80’s rigorous mathematical approaches started to be developed and, due to the great interest of the subject, they then proliferated to a multitude of different approaches. These mathematical approaches, however, are not completely equivalent as they rely on different definitions of metastable states and thus involve different properties of hitting and escape times. The situation is particularly delicate for evolutions of infinite–volume systems and for irreversible systems. The proliferation of definitions and hypotheses on metastable behavior arises from the diversity of the physical situations in which the phenomenon appears. This diversity results in contrasting demands on the appropriate mathematical theory. The main issues confronted by the mathematical treatment of metastability can be grouped into three categories.

Conservative vs non–conservative dynamics. This dichotomy applies to dynamics for statistical mechanical models of fluids or magnets. Non–conservative dynamics are those that do not conserve the total number of particles or the total magnetization. They include Glauber (that is, single spin–flip) dynamics (used to model metastable ferromagnets) and many probabilistic cellular automata (that is, parallel dynamics). In contrast, conservative dynamics are suitable to the study of supersaturated gases. Its study poses enormous challenges because particle or magnetization conservation introduces non–local effects.

Finite vs infinite configuration space. Two extreme types of metastability studies can be distinguished. In the finite–space case, the configuration space remains fixed (or bounded) while the drift towards (meta)stable configurations is increased (e.g., temperature goes to zero). In the infinite–space setup, the size of the configuration space grows in an unbounded fashion, while drifts are kept approximately constant. In many instances both parameters (drift and size of the space) change simultaneously, but usually these changes are coupled so one of the effects is dominant. Mathematically, the distinction stems from the possibility of “entropic” effects in the infinite–volume case that changes the scale at which distribution laws must be controlled. The iconic case is the thermodynamic limit of spin or gas models, see for instance [29, 30, 36, 53]. In these models, exit from metastability requires nucleation, that is the formation of a critical droplet. The probability of such an event in a large volume must include the “entropic contribution” due to the fact that the nucleation can take place.
Parallel dynamics and cost functions. Following the matrix imposed by Metropolis and Glauber dynamics, stochastic transition rates are written as exponentials of cost functions. For reversible single spin–flip dynamics these costs are determined by the difference of energy between the two states involved in the move. This is not so for parallel dynamics (for instance for Probabilistic Cellular Automata [6,20,23]) in which, at each step, all spins are independently tested. In such evolutions, costs are a possibly complicated function of the different patterns of spin flips connecting the relevant states. In these cases it is therefore necessary to dissociate energy profiles from energy barriers. The former are associated to invariant measures and determine the set of stable and metastable configurations. The latter are associated to transition rates and determine drifts and exit times.

1.2. Different approaches to metastability

Early approaches to metastability were based on the computation of expected values with respect to restricted equilibrium states [51]. This paradigm is still influential in physics, see e.g. [39]. The influence of Probability on Statistical Mechanics led to an alternative pattern of thoughts based on stochastic dynamics and focused on their spectral properties and on the behavior of their typical trajectories. This point of view has given rise to different theoretical constructions that can be classified, roughly, in three major groups.

(I) Classical approach: Hitting times of Markov chains. The escape time from metastability is determined, by the visiting or hitting time to a set of states of very small (invariant) measure, when most of the measure is carried by a different, somehow reduced, set of states (see, e.g. [37]). Similar problems were confronted in reliability theory where typical states were called good and those concerned by the hitting times were called bad. The exponential character of good–to–bad transitions is due to the existence of two different time scales: long times are needed to go from good to bad states, while the return to good states from anywhere (except, perhaps, the bad states) is much shorter. As a result, a system in a good state can reach the bad ones only through a large fluctuation bringing it all the way to the bad state. Indeed, any “intermediate” fluctuation would be followed by an unavoidable return to the good states, where, by Markovianity, the process would start afresh independently of previous attempts. Hence, the escape time is made of a large number of independent returns to the good states followed by a final successful excursion to badness that must happen without hesitation, in a much shorter time. The exit time is, therefore, a geometric random variable with extremely small success probability. In the limit exponentiality is found.

(II) Large deviations of trajectory distributions. Freidlin and Wentzel [28] were the first to use the large deviation machinery to study the problem of exit times from an attractive domain. Their theory applies to stochastic differential equations with a deterministic driving
gradient force and a small Brownian stochastic term. The deterministic part of the dynamics is responsible for the fast return to “good states” while the stochastic contribution provides the escape mechanism. The smallness of this last contribution leads to very long time scales for the visit to “bad” states. Typical trajectories are described using a graphical method built out of two basic ingredients: cycles (associated to metastable pieces of trajectories) and exit tubes (describing typical escape trajectories).

The Freidlin and Wentzel theory evolved into two related schools that we shall call the graphical [15] and the pathwise [14] approach. The former relies on a refinement of Freidlin and Wentzel’s graphical methods allowing for a detailed study of exit paths via a decomposition into cycles and saddle points traversed from cycle to cycle. The exit time also decomposes into the time spent at each point of the exit path. This graphical approach has been applied to reversible Metropolis dynamics and to simulated annealing in [15–17,55–57].

The pathwise approach, on the other hand, proposed in [14], was introduced as an adaptation of the ideas of Freidlin and Wentzel to Metropolis-like dynamics, with all notions and properties expressed in terms of an energy profile associated to the invariant measure. This provides a clearer and physically appealing picture. In particular, absolute energy minima identify stable states and “deep” local energy minima lead to metastability. The two time scales within each energy well correspond, respectively, to fast “downhill” and infrequent “uphill” trajectories. In the limit of very steep wells (temperature tending to zero), the theory yields rather precise information on: (i) the transition time, i.e., the time needed to arrive to the stable equilibrium, which is determined by the height of the largest energy barrier separating metastable from stable states. (ii) The typical exit tube, i.e., the sequences of configurations along which the formation of the stable phase takes place. This is the physically relevant mechanism that, in gas or spin systems, is mediated by the appearance of a critical droplet after which the system quickly relaxes to equilibrium.

The full power of this method was first exploited in [48] and comprehensively reviewed in [50]. It has been extended to non-reversible Markov processes in [49] (though irreversibility brings back to the graphical approach). The approach was further simplified in [40] where transition times are determined on the basis of a ranking of stability levels, without requiring detailed knowledge of typical trajectories.

(III) Potential-theoretical approach: spectral properties of Markov transition matrices. In the early eighties Aldous and Brown [1] proposed a new approach to the hitting-time theory based on spectral properties of the transition matrix and the use of the Dirichlet form. This approach has the advantage of leading to quite precise error bounds for the exponential approximation. The current version of this strategy is the potential-theoretic approach developed in [8] (see [7] for reviews). Besides exponential laws, this method gives more precise
estimates of the expected value of the transition time, including a prefactor that cannot be found with alternative approaches. The determination of this prefactor, however, requires the knowledge of the critical droplet and neighboring configurations; information that has to be imported from more detailed pathwise studies. In [3] another use of spectral and potential theoretical techniques is proposed in which only visits to well bottoms are registered. Upon time rescaling, a continuous–time Markov process is obtained whose transition rates encode the information on transition times. (See also [4,5] for recent development).

1.3. Application overview

We outline briefly some applications of the theories described above. The aim is not to be exhaustive, but rather to list references useful in relation to the definitions and comparisons to follow.

The general theory [15,16,40,48–50,54] and metastability studies in the nineties (see [2,14,19,21,26,38,43,46,47] for the pathwise approach and [17,55–57] for the graphical approach) refer to single–spin flip dynamics (Metropolis, Glauber) of Ising–like models (including mean–field versions) in finite volume and at low temperature. Studies within the general potential theoretical approach refer both to a general Markov–process point of view [9,10] and applications to mean–field and Ising model [8,13].

The study of metastability for conservative dynamics started a decade after and initially involved nearest–neighbor lattice gases at low temperature and density inside finite boxes with open boundary conditions mimicking infinite reservoirs fixing particle density. Relevant references are [31,32,36,44] for the pathwise approach and [11,33–35] for the potential theoretic approach.

Models with parallel dynamics were studied first from a numerical point of view in [6] and then rigorously in [20,22–25] (pathwise approach) and [45] (potential theoretic approach).

The more involved infinite–volume limit, at low temperature or vanishing magnetic field, was first studied via large deviations techniques in [18,27,41,42,52,53] and potential theoretically in [12]. These references dealt with Ising and Blume–Capel models under Glauber dynamics. The Ising lattice gas model subjected to Kawasaki dynamics was studied in [29,30,36] and [12] (potential theoretic approach) in the limit of temperature and volume growing exponentially fast to infinity.

1.4. Aim of the paper

As it has been explained above, due to the great theoretical and applicative interest of metastability, different mathematical theories of metastability have been developed in the past years. It is interesting to understand the mutual connections in order to apply results proven in one framework to systems that are naturally approached in a different one.
We have also remarked that within each theory different flavors have appeared. In particular, in the framework of the “large deviation” point of view, two different approaches to cycle decomposition, the graphical and the path one, have been developed. The former has been introduced and applied in a very general setup while the latter as been first introduced for Metropolis dynamics and then extended to models with cost functions such as Probabilistic Cellular Automata [20,23]. Due to its generality, the results on hitting times proven with the graphical approach are written in terms of complicated expression whose physical meaning is sometimes difficult to be caught. On the other hand, in the framework of the path approach everything is clearly written in terms of differences of energies, but the theory applies only to Metropolis–like systems.

Our opinion is that it is interesting and it can also be very fruitful to understand the connections between these two approaches in order to mix the strength of the former with the simplicity of the latter. As a first step in this direction we prove the equivalence of the two approaches in the case of the Metropolis dynamics, whose definition is now recalled.

1.5. Setup

We consider a finite state space $S$ equipped with a function $H : S \rightarrow \mathbb{R}$. Sometimes for a point $x \in S$, we will write the energy of $x$ for the value $H(x)$.

We assume that $S^2$ is equipped with a connectivity function $q : S^2 \rightarrow [0; 1]$, which satisfies the following conditions: (i) for any $x \in S$, $\sum_{y \neq x} q(x,y) \leq 1$; (ii) for any $(x,y) \in S^2$, $q(x,y) = q(y,x)$; (iii) for any $(x,y) \in S^2$, there exists $n \in \mathbb{N}$ and $x_1,\ldots,x_n$ such that $x_0 = x$, $x_n = y$ and $q(x_i,x_{i+1}) > 0$ for $i = 0,\ldots,n-1$.

For $\beta > 0$, we then define the Metropolis Markov chain $X$ as being the Markov chain with transitions given by $p_\beta$ where the kernel $p_\beta$ satisfies

$$p_\beta(x,y) = q(x,y) \exp\left(-\beta (H(y) - H(x))^+\right)$$

if $x \neq y$ and $p_\beta(x,x) = 1 - \sum_{y \neq x} p_\beta(x,y)$.

The chain started at $x \in S$ will be denoted by $x_0 = x, x_1,\ldots,x_t,\ldots$ and the associated probability will be denoted by $\mathbb{P}_x$. The main notion the paper will deal upon is that of hitting time

$$\tau_G := \inf\{t, x_t \in G\}$$

(1.1)

to a set $G \subset S$.

A particularly famous example of Metropolis dynamics is given by the standard Ising model under Glauber dynamic.

The purpose of this paper is to discuss the first exit problem of a general set $G$ for the dynamics defined above for large $\beta$. It turns out that the most relevant case corresponds
to $G$ given by a cycle, namely a set whose internal points, for large $\beta$, are typically visited many times by our process before exiting.

The paper is organized as follows: in Section 2 and Section 3 we recall the definition of cycles and the main results on hitting times respectively in the framework of the path and the graph approach. In Section 4 we explore the connections between the two approaches and state their equivalence.

2. Path cycles

We briefly review the path approach to cycle decomposition in [48]. In particular we discuss few properties that will be useful in Section 4. A path $\omega$ is a sequence $(\omega_1, \ldots, \omega_n)$ of communicating states, that is to say $q(\omega_i, \omega_{i+1}) > 0$ for any $i = 0, \ldots, n - 1$. We write $\omega : x \rightarrow y$ to denote a path joining $x$ to $y$.

We say that a subset $G$ of $S$ is connected if for any $x, x' \in G$, there exists a path $\omega : x \rightarrow x'$ such that $\omega$ is entirely contained in $G$. Two not empty subsets $G, G' \subset S$ are connected whenever there exists $x \in G$ and $x' \in G'$ such that $q(x, x') > 0$.

If $G$ is a subset of $S$ on which $H$ is constant, we will write $H(G)$ for the value of $H$ on $G$.

Definition 2.1 Let $G \subset S$, we define the exterior boundary $\partial G$ of $G$ and the ground $F(G)$ of $G$ respectively as

$$\partial G := \{ y \in S \setminus G, \exists x \in G, q(x, y) > 0 \} \quad \text{and} \quad F(G) := \{ x \in G, H(x) = \min_{G} H \}$$

For a subset $G \subset S$ and $x \in S$, we say that $x$ is a neighbor of $G$ if $x \in \partial G$.

Definition 2.2 The set $A \subset S$ is a non-trivial path cycle of $S$ if and only if it is a connected subset of $S$ verifying

$$\max_{A} H < \min_{\partial A} H$$

We say that a subset $A \subset S$ is a cycle if and only if $A$ is a singleton or $A$ is a non-trivial path cycle.

In other words, a singleton is a trivial path cycle if and only if it is not a local minimum of $H$.

Lemma 2.3 (Proposition 6.7 [50]) Given a state $x \in S$ and a real number $c \geq H(x)$, the set of all points connected to $x$ by paths whose points have energy smaller or equal to $c$ is either the trivial path cycle $\{x\}$ or is a non-trivial path cycle containing $x$. 
In the particular case $c = H(x)$ we denote by $U_{\leq x}$ the path cycle whose existence is ensured by the above lemma. In words, $U_x$ is the path cycle made of all the points in $S$ connected to $x$ via a path whose points are at energy smaller or equal to $H(x)$.

**Lemma 2.4** (Proposition 6.8 [50]) Let $A_1, A_2 \subset S$ be two path cycles such that $A_1 \cap A_2 \neq \emptyset$. Then either $A_1 \subset A_2$ or $A_2 \subset A_1$.

As a trivial consequence of these two lemmas, note that given a cycle $A$ and $x \in A$, one has the inclusion $U_{\leq x} \subset A$. On the other hand, as soon as $x \in A$ satisfies $H(x) = \max_A H$, one gets the equality $U_{\leq x} = A$.

We remark, now, the following interesting property: two non-trivial disjoint cycles cannot be connected. More precisely we state the following lemma.

**Lemma 2.5** Let $A_1, A_2 \subset S$ be two path cycles. If $A_1$ and $A_2$ are connected, then either $|A_1| = 1$ or $|A_2| = 1$.

**Proof.** By contradiction, assume that both $|A_1| > 1$ and $|A_2| > 1$ (and hence that they both satisfy equation (2.2)).

Since the two path cycles $A_1$ and $A_2$ are connected, there exists $x_1 \in A_1 \cap \partial A_2$ and, thus, it follows that

$$\max_{A_1} H \geq H(x_1) \geq \min_{\partial A_2} H > \max_{A_2} H$$

where in the last bound we used (2.2)). But similarly, there exists $x_2 \in A_2 \cap \partial A_1$, which implies

$$\max_{A_2} H \geq H(x_2) \geq \min_{\partial A_1} H > \max_{A_1} H$$

which contradicts the above inequality. \[\square\]

**Definition 2.6** Given a non-trivial path cycle $A$, we let the depth $\Gamma(A)$ and the resistance height $\tilde{\Gamma}(A)$ of the cycle be respectively

$$\Gamma(A) := \min_{\partial A} H - \min_A H = H(F(\partial A)) - H(F(A))$$

and

$$\tilde{\Gamma}(A) := \max_A H - \min_A H = \max_A H - H(F(A))$$

The following result describes the way the Markov chain exits a cycle $A$ at very low temperature. As a matter of fact, it is known in a more general, not reversible setup satisfying suitable hypotheses called the Freidlin–Wentzell conditions.
Theorem 2.7 (Theorem 6.23 in [50]) Given a non-trivial cycle $A$ and $\varepsilon > 0$, for any $x, x' \in A$, the following properties hold in the asymptotic $\beta \to \infty$:

\[ P_x \left[ \exp\{\beta(\Gamma(A) + \varepsilon)\} > \tau_{\partial A} > \exp\{\beta(\Gamma(A) - \varepsilon)\} \right] = 1 - o(1) \quad (2.3) \]

and

\[ P_x \left[ \tau_{x'} < \tau_{\partial A}, \tau_{x'} < \exp\{\beta(\tilde{\Gamma}(A) + \varepsilon)\} \right] = 1 - o(1) \quad (2.4) \]

Roughly speaking, equation (2.3) states that, starting from any point of the cycle $A$, the exit time from $A$ is of order $\exp\{\beta \Gamma(A)\}$ in the large $\beta$ asymptotic. On the other hand, starting from any point in $A$, equation (2.3) says that, before exiting $A$, the Markov chain visits all the configurations in $A$ within a time of order $\exp\{\beta \tilde{\Gamma}(A)\}$.

3. Graph cycles

The construction of graph cycles due to Freidlin Wentzell [28] is performed recursively. Here, we recall this construction following [57, Part 2]. In Section 3.2 we discuss an example.


Before starting with the recursive construction we need to recall some general definitions: given a not empty set $M$ we denote by $\mathcal{P}(M)$ the collection of all the subsets of $M$. Moreover, a function $f : M \times M \to \mathbb{R}^+ \cup \{\infty\}$, namely, a function associating each pair of elements of $M$ with a (not necessarily finite) non negative real number will be called a cost function on $M$. A path of elements of $M$ is an element $(m_1, \ldots, m_n) \in M^n$ for some $n$ positive integer.

We shall misuse the notation by also writing

\[ f(m) = \sum_{i=1}^{n-1} f(m_i, m_{i+1}) \]

for any path $m = (m_1, \ldots, m_n) \in M^n$.

The following recursive construction can be read together with the example developed in Section 3.2.

Recalling that the setup is the one introduced in Section 1.5, we define the zero-order set of graph cycles $E^0 := \{\{i\}, i \in S\}$ and the associated zero order cost function $V^0(\{i\}, \{j\}) := (H(j) - H(i))^+$ if $i$ and $j$ are connected and $V^0(\{i\}, \{j\}) := \infty$ otherwise.

It is important to remark that graph cycles [28] are usually introduced in a more general setup. For Markov chain it is usually assumed the so-called Freidlin–Wentzel assumption, namely, there exists $\kappa > 1$ such that

\[ (1/\kappa)q(x, y) \exp\{-\beta V(x, y)\} \leq p_\beta(x, y) \leq \kappa q(x, y) \exp\{-\beta V(x, y)\} \]

Note that the Metropolis dynamics is just a particular case.
Assume, then, that the $k$-order set of graph cycles $E^k \subset P(S)$ is constructed and equipped with the $k$-order cost function $V^k$. To implement the recursion, we proceed in five steps:

1. for $A \in E^k$, let
   \[ H^k_e(A) := \min\{V^k(A, A'), A' \in E^k\}. \] (3.5)
   Moreover, we define the renormalized cost function $V^k_*$ on $E^k$ by setting
   \[ V^k_*(A, B) := V^k(A, B) - H^k_e(A) \] (3.6)
   for all $A, B \in E^k$.

2. For $A, A'$ elements of $E^k$, define the $\rightarrow_k$ relation by $A \rightarrow_k A'$ if and only if there exists a path $\omega$ of elements of $E^k$ starting from $A$ and ending in $A'$ such that the cost of $\omega$ with respect to $V^k_*$ is zero, in short $V^k_*(\omega) = 0$.

3. For $A, A'$ elements of $E^k$, define the relation of equivalence $R_k$ by $A R_k A'$ if and only if both $A \rightarrow_k A'$ and $A' \rightarrow_k A$. Then we stick together all the distinct classes of equivalence and define the set $D^k+1 := \left\{ \bigcup_{A' : A R_k A'} A', A \in E^k \right\}$.

4. On the set $D^k+1$, define the (partial) order $\geq^k+1$ by $A \geq^k+1 A'$ if and only if there exist $B, B' \in E^k, B \subset A, B' \subset A'$ such that $B \rightarrow_k B'$. Then we introduce $D^k+1_\star$ the set of minimal elements for the order relation $\geq^k+1$.

5. Define $E^{k+1}$ as being the union of the set $D^k+1_\star$ and of the elements of $E^k$ which are not subsets of $D^k+1_\star$, namely
   \[ E^{k+1} := D^k+1_\star \cup \{A, A \in E^k, \exists B \in D^k+1 \setminus D^k+1_\star, A \subset B\}. \] (3.7)
   For $A \in E^{k+1}$, define
   \[ H^{k+1}_m(A) := \max\{H^k_e(A'), A' \in E^k, A' \subset A\}. \] (3.8)

6. Finally define the cost function $V^{k+1}$ on $E^{k+1}$ by
   \[ V^{k+1}(A, A') := H^{k+1}_m(A) + \min\{V^k_*(B; B'), B, B' \in E^k, B \subset A, B' \subset A'\}. \] (3.9)

The construction continues until $E^k = \{S\}$. As noted in [57], the recursive procedure described here is not stationary until iteration $n_S$, where $n_S$ is the first iteration such that $E^{n_S} = S$. 

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We remark that for any $k \geq 0$, $E^k$ is a partition of $S$ and more precisely the procedure gives a hierarchical decomposition of the state space as a tree starting from the singletons and ending with the whole space.

We define the set of graph–cycles $\mathcal{C} := \bigcup_{k \geq 0} E^k$ and call any element of $\mathcal{C}$ a graph–cycle.

**Definition 3.8** Let $A \in \mathcal{C}$. We set

$$H_e(A) := \sup_{k \geq 0} \{H^k_e(A)\}$$

if $A \neq S$ and $H_e(S) = \infty$.

The above definition is based on the following remark (see [57]): for $A \in \mathcal{C}$, it is easy to see that, whenever $A \in E^k \cap E^{k+1}$, one has

$$H^{k+1}_m(A) = H^{k+1}_e(A) = H^k_e(A)$$

and thus $H^k_e(A) = H_e(A)$ as soon as $A \in E^k$.

**Definition 3.9** Let $A \in \mathcal{C}$ such that $A \neq S$ and $|A| > 1$. We set

- $\mathcal{C}^*_A := \{B \in \mathcal{C}, B \subset A, B \neq A\}$;
- $\mathcal{M}_*(A) := \{B \in \mathcal{C}, B$ is a maximal element in $\mathcal{C}^*_A\}$ (maximal proper partition of $A$);
- $H_m(A) := \sup\{H_e(B), B \in \mathcal{C}^*_A\} \lor 0$.

We now state the analogous of Theorem 2.7 in the framework of graph–cycles. The proof is given, for instance, in [16].

**Theorem 3.10** (Propositions 4.19, 4.20 and 5.1 [16]) Let $A$ be a graph cycle. For any $\varepsilon > 0$, for any $x, x' \in A$, as $\beta \to \infty$, one has the asymptotic:

$$\mathbb{P}_x\left[\exp\{\beta(H_e(A) + \varepsilon)\} > \tau_{\partial A} > \exp\{\beta(H_e(A) - \varepsilon)\}\right] = 1 - o(1)$$

and

$$\mathbb{P}_x[\tau_{x'} < \tau_{\partial A}, \tau_{x'} < \exp\{\beta(H_m(A) + \varepsilon)\}] = 1 - o(1).$$

We note that this result strongly suggests the equalities $H_m(A) = \Gamma(A)$ and $H_e(A) = \tilde{\Gamma}(A)$. In the next section we shall prove that this fact is indeed true.
Figure 3.1: The energy landscape on $S = \{a, b, c, d, e, f, g, h, i, j, k\}$

3.2. An example.

In this part, we run the algorithm described above in a simple case. We consider the state space $S = \{a, b, c, d, e, f, g, h, i, j, k\}$ with connectivities and energy landscape described in Figure 3.1. For example, $H(a) = H(d) = 2, H(i) = 0, q(d, e) > 0, q(e, f) > 0$ and $q(e, g) = 0$.

Iteration 1. To start our construction, we first recall that

$$E^0 = \{\{a\}, \{b\}, \{c\}, \{d\}, \{e\}, \{f\}, \{g\}, \{h\}, \{i\}, \{j\}, \{k\}\}.$$

For two singletons which are not connected (say for example $a$ and $c$), we then have by definition the equality $V^0(\{a\}, \{c\}) = \infty$. Else, it is easy to see that $V^0(\cdot, \cdot)$ is equal to zero for connected singleton except in the following cases:

$$V^0(\{a\}, \{b\}) = 3, V^0(\{c\}, \{d\}) = 4, V^0(\{c\}, \{d\}) = 1, V^0(\{f\}, \{g\}) = 2,$$

$$V^0(\{h\}, \{g\}) = 1, V^0(\{i\}, \{j\}) = 1, V^0(\{j\}, \{h\}) = 3, \text{ and } V^0(\{j\}, \{k\}) = 4.$$

One can then compute the following quantities using (3.5)

$$H^0_e(\{a\}) = 3, \ H^0_e(\{c\}) = 1, \ H^0_e(\{i\}) = 1,$$

and

$$H^0_e(\{b\}) = H^0_e(\{d\}) = H^0_e(\{e\}) = H^0_e(\{f\}) = H^0_e(\{g\}) = H^0_e(\{h\}) = H^0_e(\{j\}) = H^0_e(\{k\}) = 0.$$
By making use of equation (3.10) and of the remark after it, we note that, for every \( x \in S \) and for all \( k \geq 0 \), \( H_e^k(\{x\}) = H_e^0(\{x\}) = H_e(\{x\}) \). In other words the remark above gives the value of the functions \( H_e \) and \( H_e^k \), for \( k \geq 1 \), computed at any singleton, namely \( H_e(\{x\}) \) and \( H_e^k(\{x\}) \) for \( k \geq 1 \) and \( x \in S \). A similar observation will apply in the next steps of the algorithm.

Using (3.6), we now compute \( V_0^0(\cdot, \cdot) \), which is infinite for not connected singletons and zero for all connected singletons apart from the following cases:

\[
V_0^0(\{c\}, \{b\}) = 3, V_0^0(\{f\}, \{g\}) = 2, V_0^0(\{h\}, \{g\}) = 1, V_0^0(\{i\}, \{h\}) = 2, V_0^0(\{j\}, \{k\}) = 4.
\]

Applying the third step of the construction of Section 3.1, we then deduce that

\[
D^1 = \{\{a, b\}, \{c, d, e, f\}, \{g\}, \{h\}, \{i, j\}, \{k\}\}.
\]

For example, to see that \( \{a, b\} \) is a subset of \( D^1 \), one notes that \( \{a\} \mathcal{R}_0 \{b\} \) since \( V_*^0(\{a\}, \{b\}) = V_*^0(\{b\}, \{a\}) = 0 \) and \( \{b\} \) is not in relation for \( \mathcal{R}_0 \) with \( \{c\} \) since \( V_*^0(\{c\}, \{b\}) = 3 \neq 0 \).

The set \( \{a, b\} \) is not a minimal element for \( \geq^1 \) because \( \{b\} \not\rightarrow \{c\} \), which implies that \( \{a, b\} \geq^1 \{c, d, e, f\} \). On the other hand, one checks readily that both \( \{c, d, e, f\} \) and \( \{i, j\} \) are minimal elements for \( \geq^1 \), and we deduce:

\[
D^*_1 = \{\{c, d, e, f\}, \{i, j\}\}.
\]

Combining (3.7) and (3.14), we obtain:

\[
E^1 = \{\{a\}, \{b\}, \{c, d, e, f\}, \{g\}, \{h\}, \{i, j\}, \{k\}\}.
\]

We then get immediately by (3.8) : \( H_m^1(\{x\}) = H_e(\{x\}) = 0 \) for \( x \in \{b, g, h, k\} \) and

\[
H_m^1(\{a\}) = 3, H_m^1(\{c, d, e, f\}) = 1, H_m^1(\{i, j\}) = 1.
\]

We can now compute \( V^1 \) making use of (3.9), (3.15) and (3.13). Once again, \( V^1(A, B) \) is infinite as soon as \( A \) and \( B \) are disconnected, and \( V^1(A, B) = V^0(A, B) \) as soon as \( A \in E^0, B \in E^0 \). For the remaining cases, the cycle pairs with not zero \( V^1 \) are:

\[
V^1(\{c, d, e, f\}, \{b\}) = 4, V^1(\{c, d, e, f\}, \{g\}) = 3, V^1(\{i, j\}, \{h\}) = 3, V^1(\{i, j\}, \{k\}) = 5,
\]

\[
V^1(\{b\}, \{c, d, e, f\}) = V^1(\{g\}, \{c, d, e, f\}) = V^1(\{h\}, \{i, j\}) = V^1(\{k\}, \{i, j\}) = 0.
\]

**Iteration 2.** For the second iteration of the algorithm, one proceeds as in Iteration 1 replacing the state space \( E^0 \) equipped with the cost function \( V^0 \) by the space \( E^1 \) equipped with the cost function \( V^1 \). Applying again (3.5) and the remark below (3.10), we compute:

\[
H_e(\{i, j\}) = H_e^1(\{i, j\}) = 3, H_e(\{c, d, e, f\}) = 3.
\]
Once again, $V^1_*(A, B)$ is infinite as soon as $A$ and $B$ are not connected, identically null otherwise except in the following cases:

$$V^1_*(\{c, d, e, f\}, \{b\}) = 1, V^1_*(\{i, j\}, \{k\}) = 2, V^1_*(\{h\}, \{g\}) = 1,$$

and hence we deduce

$$D^2 = \{\{a, b\}, \{c, d, e, f, g\}, \{h, i, j\}, \{k\}\}.$$ 

Since neither $\{a, b\}$ nor $\{c, d, e, f, g\}$ are minimal elements for the order relation $\geq^2$, we obtain

$$E^2 = \{\{a\}, \{b\}, \{c, d, e, f\}, \{g\}, \{h, i, j\}, \{k\}\}.$$ 

Finally we get

$$H_m(\{h, i, j\}) = 3.$$

Iteration 3. Similarly, for the third iteration of the algorithm, one shows that:

$$H_e(\{h, i, j\}) = 4,$$

$$D^3 = \{\{a, b\}, \{c, d, e, f, g, h, i, j\}, \{k\}\},$$

$$E^3 = \{\{a\}, \{b\}, \{c, d, e, f, g, h, i, j\}, \{k\}\},$$

$$H_m(\{c, d, e, f, g, h, i, j\}) = 4.$$

Iteration 4. After the fourth iteration, the procedure is complete (that is $n_S = 4$), and we get:

$$H_e(\{c, d, e, f, g, h, i, j\}) = 5,$$

$$D^4 = E^4 = S,$$

$$H_m(S) = 5, H_e(S) = \infty.$$

We finally deduce the set of graph-cycles of $S$:

$$\mathcal{C}(S) = \{\{a\}, \{b\}, \{c\}, \{d\}, \{e\}, \{f\}, \{g\}, \{h\}, \{i\}, \{j\}, \{k\},$$

$$\{c, d, e, f\}, \{i, j\}, \{h, i, j\}, \{c, d, e, f, g, h, i, j\}, S\}.$$ 

We remark on this toy example that the set of graph cycles coincides with the set of path cycles. We show in the next section that this is indeed always the case.

4. Equivalence of definitions

In this section we prove that path and graph–cycle decompositions are equivalent. As a byproduct of the proof we shall also give a physical interpretation, in terms of differences of energy, of the quantities $H_e$ and $H_m$ defined in the framework of the graph theory.
Theorem 4.11 With the notations introduced above:

1. any graph–cycle \( A \in \mathcal{C} \) is a path–cycle. Furthermore the following equality holds:

\[
H_e(A) = \left( \min_{\partial A} H - \min_A H \right) \vee 0.
\] (4.16)

One also has

\[
H_m(A) = \max_A H - \min_A H
\] (4.17)

when \( |A| > 1 \) and \( H_m(A) = H_e(A) \) when \( |A| = 1 \).

2. Any path–cycle is a graph–cycle.

In particular, starting from anywhere inside a cycle, the exit time of the cycle on an exponential scale is essentially given by the depth of the cycle. Note that as soon as \( A \in \mathcal{C} \) is a non trivial path cycle, then \( H_e(A) = \min_{\partial A} H - \min_A H \). More precisely, we have the equivalence between Theorems 2.7 and 3.10.

In order to proof item 1 of the theorem we shall proceed recursively. In particular we shall introduce the condition \( \mathcal{H}_n \) and prove that it holds for any \( n \geq 0 \).

**Condition \( \mathcal{H}_n \).** For any \( A \in E^n \) the following properties hold true:

\[
\begin{cases}
(I)_n & A \text{ is a path cycle}; \\
(II)_n & \text{if } A \text{ (where } |A| > 1 \text{) and } \{a\} \text{ are connected and belong to } E^n \text{, then } V^n(A, \{a\}) = H(a) - \min_A H \text{ and } V^n(\{a\}, A) = 0; \\
(III)_n & H^n_e(A) = (\min_{\partial A} H - \min_A H) \vee 0; \\
(IV)_n & H^n_m(A) = \max_A H - \min_A H \text{ when } |A| > 1.
\end{cases}
\]

Before turning to the proof of the theorem, we make a few remarks which hold provided \( \mathcal{H}_k \) is true for any \( k = 0, 1, \ldots, n - 1 \), and that will be used throughout the proof.

R1. One important point in the proof will be to compute \( V^n(A, B) \) for any element \( A, B \in E^n \). Note that if both \(|A| > 1 \) and \(|B| > 1 \), then it follows from Lemma 2.5 that \( A \) and \( B \) are not connected, and from (3.9) follows that \( V^n(A, B) = \infty \). On the other hand, if two singletons \( \{a\} \) and \( \{b\} \) belonging to \( E^n \) are connected, then it follows immediately from the definition of \( V^n \) that \( V^n(\{a\}, \{b\}) = (H(b) - H(a))^+ \). The other cases are covered by assumption \( (II)_n \).

R2. We state that

\[
E^{n+1} \setminus E^n = D^{n+1}_n.
\] (4.18)

Indeed, given the definition (3.7) of \( E^{n+1} \), for (4.18) to hold, we have to show that an element \( B \in D^{n+1}_n \) cannot belong to \( E^n \). Assume by contradiction that \( B \in E^n \); there exists at least
one element $B' \in E^n$ connected to $B$ such that $B \xrightarrow{n} B'$ (and hence $B \geq^n B'$). Since $B \in D_r^{n+1}$, it is a minimal element for $\geq^n$, and hence one gets $B' \geq^n B$. This implies that both $B$ and $B'$ are in the same equivalence class for $\cal{R}_n$, which is contradictory in view of the construction of $D_r^{n+1}$.

This remark is very useful, indeed, to prove the implication $\cal{H}_n \Rightarrow \cal{H}_{n+1}$, we show the four properties defining $\cal{H}_{n+1}$ restricting the analysis to the case $A \in D_r^{n+1}$. We write $A$ as the disjoint union

$$A = \bigsqcup_{j=1}^{j_A} A_j \quad (4.19)$$

where $j_A \geq 1$ and, for all $j \leq j_A$, $A_j \in E^n$. This decomposition is in fact the decomposition $\cal{M}_s(A)$ which we rewrite in a more tractable way.

Of course $A$ not belonging to $E^n$ implies in fact that $j_A \geq 2$. In the remaining of the proof, we will refer to the elements appearing in the decomposition of the right hand side of (4.19) as subcycles of $A$.

For $j \in [1, j_A]$, whenever there exists $a \in S$ such that $A_j = \{a\}$ and $b \in S$ a neighbor of $a$ such that $H(b) \leq H(a)$, we say that the singleton $A_j$ is a good singleton of $A$. In words, a good singleton $\{a\}$ of $A$ is a trivial path cycle because $a$ is not a local minimum of $H(\cdot)$. We denote by $j_A^s$ the number of subcycles of $A$ which are good singletons of $A$. Note that whereas the notion of trivial path cycles depends only on the energy landscape, the notion of good singleton of $A$ strongly depends on the cycle decomposition of $A$.

R4. $(I)_n$ implies that, as soon as $A_j \xrightarrow{n} \{a\}$, where $|A_j| > 1$ and $\{a\}$ is a subcycle of $A$ such that $a \in \partial A_j$, then necessarily $\{a\}$ is good.

R5. For singletons which are not good, say that $A_j = \{a\}$ where $j > j_A^s$, we can remark that all the subcycles of $A$ which are connected to $\{a\}$ are good singletons and have the same energy. Indeed, all these subcycles are path cycles by $(I)_n$. Since a non–trivial path cycle cannot be connected to a local minimum for $H(\cdot)$ in view of (2.2), it follows that all these subcycles are good singletons. To see that they have the same energy, consider any $b \in \partial\{a\} \cap A$ satisfying $\{a\} \xrightarrow{n} \{b\}$. By the definition of $\xrightarrow{n}$, we then get that $H(b) = H(a) + H_e(\{a\})$.

R6 Combining R4, R5 and Lemma 2.5, we note the important fact that $j_A^s \geq 1$. Up to reordering, we assume from now on that for $1 \leq j \leq j_A^s$, $A_j$ is a good singleton.

R7. We then show the following equalities:
\[
H \left( \bigcup_{1 \leq j \leq j_A^*} A_j \right) = \max_A H 
\]
(4.20)

and
\[
H (\partial A_j \cap A, j \in [j_A^* + 1, j_A]) = \max_A H.
\]
(4.21)

Of course, equations (4.20) and (4.21) implicitly state that \( H(A_j) \) does not depend on \( j \in [1, j_A^*] \) and that \( H (\partial A_j \cap A) \) does not depend on \( j \in [j_A^* + 1, j_A] \).

To prove (4.20) and (4.21), we distinguish the cases \( j_A^* = 1 \) and \( j_A^* \geq 2 \).

If \( j_A^* = 1 \), then for every \( j \in [2, j_A] \), it follows from Lemma 2 and remark R4 that \( A_j \) is connected to \( A_1 \) and \( A_j \to A_1 \) for \( j \in [2, j_A] \). As a consequence, considering \((II)_n\), we get both (4.20) and (4.21).

Assume now that \( j_A^* \geq 2 \) and consider two elements \( A_i \) and \( A_j \) (where \( 1 \leq i, j \leq j_A^* \) and \( i \neq j \)). We consider a path \( \omega = (\omega_1, \ldots, \omega_n) \) of connected subcycles of \( A \) such that for \( \omega_1 = A_i, \omega_n = A_j \) and for any \( k \in [1, n-1], \omega_k \to \omega_{k+1} \).

Let us consider \( k_0 \in [1, n-2] \) such that \( \omega_{k_0} = A_l \) where \( l \leq j_A^* \), or in words \( \omega_{k_0} \) is a good singleton of \( A \). We claim that either \( \omega_{k_0+1} \) or \( \omega_{k_0+2} \) is a good singleton of \( A \), and moreover that in both cases the energy of this good singleton is equal to the energy of \( \omega_{k_0} \).

Indeed, if \( \omega_{k_0+1} \) is not a good singleton, then either \( |\omega_{k_0+1}| > 1 \) or \( \omega_{k_0+1} \) is a local minimum for \( H \). When \( |\omega_{k_0+1}| > 1 \), it follows from remark R4 that necessarily \( \omega_{k_0+2} \) is a good singleton. Also, it follows from \((II)_n\) that \( H(\omega_{k_0}) = H(\omega_{k_0+2}) \). On the other hand, in the case where \( \omega_{k_0+1} \) is a local minimum for \( H \), we saw in remark R5 that necessarily \( \omega_{k_0+2} \) is a singleton and that its energy is the same as the one of \( \omega_{k_0} \).

It follows easily from these considerations that \( H(A_i) = H(A_j) \), and hence \( H(A_j) \) does not depend on \( j \in [1, j_A^*] \).

The equalities (4.20) and (4.21) then follow readily from Lemma 2.5, from \((II)_n\) and \((I)_n\).

**Proof of Theorem 4.11.** Proof of item 1: we shall prove recursively that Condition \( \mathcal{H}_n \) holds for any \( n \geq 0 \), which implies item 1. Since singletons are path–cycles, \( \mathcal{H}_0 \) is trivially true. Note in particular that for any \( x \in S \), the equality \( H_m(\{x\}) = H_e(\{x\}) \) holds true.

Let us now assume that \( \mathcal{H}_n \) holds for a given \( n \geq 0 \) and prove that \( \mathcal{H}_{n+1} \) holds. Considering R2, we restrict ourselves to the case where \( A \in D^{n+1}_n \) (and in particular \(|A| \geq 1 \)); in the remaining of the proof of Theorem 4.11, \( A \) will be such an element.

\((IV)_{n+1}\) is a consequence of (4.20), \((III)_n\) and \((II)_n\). Indeed, since \(|A| \geq 1 \), combining the definition (3.8) of \( H_m \) and the induction hypothesis, we get the identity
\[
H_m^{n+1}(A) = \max \left\{ \left( \min_{A'} H - \min_{A'} H \right) \right\} \vee 0, A' \in E^n, A' \subset A .
\]
(4.22)
In the case where \( j_A^* = j_A \), then (4.22) immediately gives that \( H_m(A) = 0 \), which implies (\( IV \))_{n+1} using (4.20).

Assume now that \( j_A > j_A^* \). We note that one can restrict the optimization appearing in the right hand side of (4.22) to subcycles \( A' \) of \( A \) verifying \( A' = A_j \) for \( j \in [j_A^* + 1; j_A] \); indeed, it follows from \( (I) \) \( n \) (in the case where \( |A_j| > 1 \)) and from remark R5 (when \( |A_j| = 1 \)) that the quantity \( \min_{\partial A'} H - \min_A H \) is positive for such cycles and it is identically null as soon as \( A' \) is a good singleton.

For such an \( A' \), the equality (4.21) implies that the quantity \( \min_{\partial A'} H \) does not depend on \( A' \) and is equal to \( \max_A H \). Hence one gets

\[
H_m^{n+1}(A) = \max_A H - \min_{j \in [j_A^* + 1; j_A]} \min_{A_j} H.
\]

Since \( E' \) is a partition of \( S \) and since \( \max_A H \) is reached on \( A \setminus \bigcup_{1 \leq j \leq j_A^*} A_j \) as shown in (4.20), one gets \( \min_{j \in [j_A^* + 1; j_A]} \min_{A_j} H = \min_A H \), and hence the equality (IV)_{n+1} holds in this case as well.

(III)_{n+1} is a direct consequence of (II) \( n \) and (I) \( n \). By Lemma 2.5, we get that all the neighbors of \( A \) in \( E' \) are singletons, and hence we have:

\[
H_e^{n+1}(A) = \inf_{\{a\} \in E', a \in \partial A} V^n(A, \{a\}) = \inf_{\{a\} \in E', a \in \partial A} H(a) - \min_{A_j} H
\]

where in the second equality we made use of (II) \( n \).

(IV)_{n+1} is a consequence of (III)_{n+1} and (IV)_{n+1}. Indeed, we have to show that \( \max_A H < H(F(\partial A)) \). Note that this is equivalent to showing that \( H_m^{n+1}(A) < H_e^{n+1}(A) \), which in turn by definition of \( H_e(\cdot) \) is equivalent to the fact that, for every \( A' \in E' \):

\[ \min\{V'_e(B, B'), B, B' \in E', B \subset A, B' \subset A' \} > 0. \]

By contradiction, assume that there exist \( A' \in E' \), \( A' \neq A, B, B' \in E', B \subset A, B' \subset A' \) such that \( V'_e(B, B') = 0 \) (and of course \( A \neq A' \)); this implies that \( A \geq^{n+1} A' \), which contradicts the minimality of \( A \) for the order relation \( \geq^{n+1} \) (given the construction of \( A \), this minimality is necessarily strict).

(II)_{n+1} is a consequence of (IV)_{n+1}, (III)_{n+1} and (II) \( n \). Indeed, for \( a \in \partial A \) such that \( \{a\} \in E^{n+1} \), using the definition (3.9) of \( V_n \), of (IV)_{n+1} and of (I) \( n \), we get that:

\[
V^{n+1}(A, \{a\}) = \max_A H - \min_A H + \min\{V'_e(B, \{a\}), B \in E', B \subset A\}. \tag{4.23}
\]

Note that we restricted the set on which we minimize \( V'_e \) in the right hand side of (4.23) since in all other cases this quantity is infinite.
Now we distinguish two cases.

In the case where the min in the right hand side of (4.23) is attained on a subcycle \( A_{j_0} \) of \( A \) such that \( j_0 > j_A \), making use of \((II)_n\) and \((III)_n\), we get:

\[
\min \{ V^n(B, \{a\}), B \in \mathcal{E}^n, B \subset A \} = H(a) - \min_{\partial A_{j_0}} H - (\min_{\partial A_{j_0}} H - \min \{ V^n(B, \{a\}), B \in \mathcal{E}^n, B \subset A \}) = H(a) - \min \{ V^n(B, \{a\}), B \in \mathcal{E}^n, B \subset A \}. \tag{4.24}
\]

Then we already noted that \( \min_{\partial A_{j_0}} H = \max_A H \) holds because \( j_0 > j_A \) and of equality (4.21).

When the min in the right hand side of (4.23) is attained on a good singleton \( A_{j_1} \) with \( j_1 \in [1, j^*_A] \), one gets that

\[
\min \{ V^n(B, \{a\}), B \in \mathcal{E}^n, B \subset A \} = (H(a) - H(A_{j_1}))^+ = H(a) - \max_A H. \tag{4.25}
\]

Equation (4.25) holds because \( A \) satisfies (2.2), and hence \((H(a) - H(A_{j_1}))^+ = H(a) - H(A_{j_1})\), and on the other hand \( H(A_{j_1}) = \max_A H \) from (4.20).

Finally, we have that \( V^{n+1}(\{a\}, A) = H_m(\{a\}) + \min \{ V^n(\{a\}, A_j) - H_e(\{a\}), j \in [1, j_A] \} = \min \{ V^n(\{a\}, A_j), j \in [1, j_A] \} \). When this minimum is attained on a non–trivial subcycle of \( A \), one gets \( V^{n+1}(\{a\}, A) = 0 \) by \((II)_n\), and when it is attained on a singleton \( \{b_0\} \) involved in the decomposition (4.19), using \((I)_{n+1}\), one gets that \( V^{n+1}(\{a\}, A) = ((H(b_0) - H(a))^+ = 0 \), thus in any case

\[
V^{n+1}(\{a\}, A) = 0. \tag{4.26}
\]

Combining (4.23), (4.25), (4.24) and (4.26), we get \((II)_{n+1}\), and hence the recursion is completed. As noted above this completes the proof of item 1.

Now we prove item 2 of the theorem. Let us consider \( A \) a connected subset of \( S \) which is not a singleton, and such that \( \max_A H < \min_{\partial A} H \).

We introduce \( k_0 := \inf \{k \geq 1, \exists A' \in \mathcal{E}^k, A \subset A' \} \) and we denote by \( A' \) the element of \( \mathcal{E}^{k_0} \) such that \( A \subset A' \). Since the recursion of section 3 is not stationary, \( k_0 \) is well defined. To show Theorem 4.11, it is enough to show that the reciprocal inclusion holds. We introduce the decomposition of \( A' \) into disjoint elements of \( \mathcal{E}^{k_0-1} \) as we did in (4.20):

\[
A' = \bigcup_{j=1}^{j_{A'}} A'_{j}, \tag{4.27}
\]

and we define \( j_A' \) for \( A' \) as we defined \( j_A^* \) for \( A \) in R3.

The inclusion \( A' \subset A \) will hold as soon as we show that there exists an element \( a \in S \) such that \( \{a\} \) is a good singleton for \( A' \) and \( a \in A \); indeed, we already proved in item 1 that \( A' \) is a path cycle, hence combining Lemma 2.3 and equality (4.20), we remark that \( A' = U_{\leq a} \). From Lemma 2.4, we thus get \( A' \subset A \).
By the minimality of $k_0$, $A$ is not contained in any of the $(A'_j)_{j \leq j'}$, and thus there exist $j_1$ and $j_2$ in $[1; j']$ such that both $A \cap A'_{j_1} \neq \emptyset$ and $A \cap A'_{j_2} \neq \emptyset$. $A$ being connected, there exists a path $\omega = (\omega_1, \ldots, \omega_n)$ contained in $A$ joining $A \cap A'_{j_1}$ to $A \cap A'_{j_2}$ (say $\omega_1 \in A_{j_1}$ and $\omega_n \in A_{j_2}$); we then consider three cases:

1. either $A'_{j_1}$ or $A'_{j_2}$ is a good singleton, and then we are done.

2. either $A_{j_1}$ or $A_{j_2}$ is not a singleton, and then we assume by symmetry that $|A_{j_1}| > 1$.

Let us consider $l = \inf_{j \geq 1} \{\omega_j / \omega_j \notin A_{j_1}\}$; by Lemma 2.5, $\{\omega_l\}$ is necessarily a subcycle of $A'$, and by R4, $\{\omega_l\}$ is a good singleton.

3. in the case where both $A_{j_1}$ and $A_{j_2}$ are not good singletons, then it follows from R5 that they cannot be connected and hence $n \geq 3$. It follows from R5 as well that $\{\omega_2\}$ is necessarily a good singleton.

□

References


